

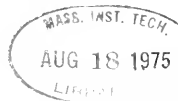
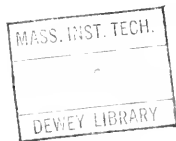
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AN ISO-CONTOUR PLOTTING ROUTINE AS A TOOL
FOR MAXIMUM LIKELIHOOD ESTIMATION

RAFAEL C. ANDREU

June 1975

WP 795-75

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doctoral program at the MIT Sloan School
of Management.

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1.- Introduction.

Doing maximum likelihood estimation when the likelihood function is so complicated that it becomes very difficult to deal with it analytically is a problem which has to be solved numerically.

The process of estimating the parameters of an underlying lognormal distribution when sampling is considered to be without replacement and proportional to size constitutes one such case. This particular sampling process is being used as a model for inferring the distribution of yet undiscovered oil and gas deposits from a sample of discovered pools. Pool discovery (observing deposit sizes in order of discovery) "is more akin to sampling without replacement and proportional to size than to sampling values of independent, identically distributed random variables" (1). This feature had been ignored until recently.

In addition, as it is also pointed out in (1), the population from which the discovery process picks pool sizes can be considered as a sequence of N values of N mutually independent random variables identically distributed with a common lognormal density, generated by nature. There is empirical evidence showing that a variety of actual size distributions can be reasonably characterized as being lognormal.

The mathematical analysis of such a sampling process is contained in (1), where an asymptotic expansion of the likelihood function (valid for large N and fixed sample

size n) is developed, since working with the exact expressions turns out to be very difficult.

In turn, since this asymptotic expansion is analytically complicated, to understand its behavior we supplement the treatment of it in (1) with numerical analysis of particular cases. Its complexity may be better understood by looking at its expression. The expansion for the sampling density of \underline{Y} , valid for large $p = N-n$ is:

$$I_{N,n}(\underline{Y}) \prod_{j=1}^n Y_j f(Y_j | \underline{\theta}) = \left\{ \frac{\Gamma(p+n+1)}{\Gamma(p+1)} \prod_{j=1}^n \frac{Y_j f(Y_j | \underline{\theta})}{[pM_1 + b_j]} \right\} \\ \times \left\{ 1 + \frac{1}{2} p V g_2(pM_1, \underline{Y}) + \frac{1}{8} p^2 V^2 g_4(pM_1, \underline{Y}) \right. \\ \left. + p \left(\frac{1}{6} M_3 - \frac{1}{2} M_1 M_2 + \frac{1}{3} M_1^3 \right) g_3(pM_1, \underline{Y}) + O(p^{-3}) \right\}$$

where the functions $g_m(pM_1, \underline{Y})$ are given by

$$g_2(pM_1, \underline{Y}) = \sum_{j=1}^n (pM_1 + b_j)^{-2} + \left[\sum_{j=1}^n (pM_1 + b_j)^{-1} \right]^2, \\ g_3(pM_1, \underline{Y}) = \left[\sum_{j=1}^n (pM_1 + b_j)^{-1} \right]^3 + 3 \left[\sum_{j=1}^n (pM_1 + b_j)^{-1} \right] \left[\sum_{j=1}^n (pM_1 + b_j)^{-2} \right] \\ + 2 \sum_{j=1}^n (pM_1 + b_j)^{-3},$$

$$\begin{aligned}
g_4(pM_1, \underline{Y}) = & \left[\sum_{j=1}^n (pM_1 + b_j)^{-1} \right] g_3(pM_1, \underline{Y}) \\
& + 3 \left[\sum_{j=1}^n (pM_1 + b_j)^{-1} \right]^2 \left[\sum_{j=1}^n (pM_1 + b_j)^{-2} \right] \\
& + 3 \left[\sum_{j=1}^n (pM_1 + b_j)^{-2} \right]^2 + 6 \sum_{j=1}^n (pM_1 + b_j)^{-4} \\
& + 6 \left[\sum_{j=1}^n (pM_1 + b_j)^{-1} \right] \left[\sum_{j=1}^n (pM_1 + b_j)^{-3} \right] .
\end{aligned}$$

and where \underline{Y} is a vector of random observations, $f(Y | \underline{\theta})$ is the lognormal density with vector parameter $\underline{\theta} = (\mu, \sigma^2)$, M_1 and M_2 are the first two moments of such density, N = population size, n = sample size and $b_j = \sum_{i=1}^j Y_i$.

Our main goal in studying this expansion is to obtain approximate maximum likelihood estimates (MLE) for the parameters of the underlying lognormal density, μ and σ^2 , and for the finite population size N . Namely, defining the likelihood function for μ, σ^2 and N given a data vector \underline{Y} as $l(\mu, \sigma^2, N | \underline{Y})$, we wish to find a triple (μ, σ^2, N_0) of values (μ, σ^2, N) that maximize l .

ML estimators may be analytically obtained in the limit, as $N \rightarrow \infty$; hence a question of interest is how large should N be to actually obtain in practice a MLE of all three parameters μ, σ^2 and N . It is also interesting to get an idea of how stable these estimates turn out to be as a function of the sample size.

The approach we took to numerically analyze the behavior of the likelihood function for particular samples was to construct iso - contour graphs of it. Doing so in two dimensions with one parameter fixed allowed us to visually de-

tect interesting properties of the function.

What we did was the following: Taking the population size N as fixed and regarding μ and σ^2 as arguments of the likelihood function, we generated sets of points in the (μ, σ^2) plane satisfying

$$l(\mu, \sigma^2, N | \underline{y}) = k,$$

with k fixed. (Such sets of points are called iso - contours). Plotting sets of such iso - contours for different values of the population size and different values of the sample size, we gained some insight into properties of the likelihood function.

In what follows we will describe these properties , along with the estimates we obtained for the parameters mentioned above.

In a different context, certain methodological details of the design of an iso - contour generating routine turn out to be interesting as well. While it is easy to find a set of points satisfying an iso - contour equation (over given intervals for the function's arguments, and provided we have tabulated function values for points belonging to those intervals), what constitutes a more difficult problem is to order them in such a manner that a plotting device can go from one point to the next actually drawing the iso - contour curve. Thus, we will also describe the method we used in writing the routine which generates ordered sets of points belonging to a given contour. Some details are particularly relevant when one is interested in the function's extreme

points; these will be emphasized in the corresponding section. Program details, including actual program listings, are included in Appendix A.

As for the routine implementation, we have it working in the TROLL system environment, taking advantage of certain plotting facilities already available in this system.

A couple of macros which make the routine invocation easier from TROLL are included in Appendix B, along with their operating procedures.

Finally, some data referring to program requirements (memory and execution times) are outlined in Appendix C.

2.- Iso - contour generating routine. Outline of the method employed.

As stated in the previous section, the main difficulty lies on the proper ordering of the set of points belonging to a given iso - contour. The procedure we followed, inspired by that described in (2), orders the sets of points as it generates them and solves other problems that sometimes arise, such as having two or more disjoint iso - contour branches in the intervals of interest for the function's arguments.

Although some published material was available describing programs also motivated by the method introduced in (2) ref. (3), (4) , they were badly documented and not particularly well suited to interface with TROLL. Further, some interesting problems arise when generating an iso - contour near the function's extreme points, which were precisely the ones we were more interested in. Our final procedure differs slightly from that in (2) regarding to these problems, and so we thought it was worthwhile to briefly describe it here.

A general outline of the method will be presented first, so that those problems may be more easily understood.

What the routine accepts as input data is a table of function values in the points of a rectangular grid defined by a set of rectangular intervals of the function's arguments, with increments that may vary along the intervals. In such a grid, two units turn out to be of importance when

the aim is to generate iso - contours; namely the grid edges and the grid cells.

Grid edges are used as units during the first step of the procedure, in which each one is checked to see whether or not it is crossed by the iso - contour being built. A way to conduct such a check is to test the condition

$$(f_1 - k) \cdot (f_2 - k) \leq 0 \quad (a)$$

for every edge in the grid, where f_1 and f_2 are the function's values at the edge extreme points and k is the level of the desired iso - contour. Successful tests are recorded during this first step, so that when it is finished (having tested all the grid edges) , we are left with a set of grid edges crossed by the iso - contour under construction.

Two properties of this procedure must be emphasized at this point. First, test (a) will only detect at most one crossing point at any edge (so that when the edges' size is too big, situations such as the one depicted in Fig. 2.1 will not be properly recorded), and points of tangency

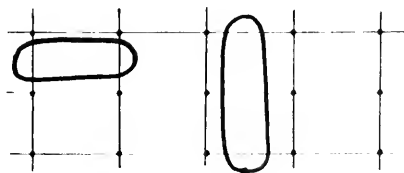


Fig. 2.1

between grid edges and iso - contour lines will not be recorded at all. Second, the procedure will go on in its second

step to locate the points of intersection between edges and contour lines, so that the points contained in the resulting iso - contour set will all lie on grid edges.

The second step of the procedure takes grid cells as working units, and relies on the fact that when the grid size is properly chosen, either none or two edges of a grid cell will be crossed by an iso - contour line (in general, if tangency points are not recorded, an even number of edges will be crossed in any grid cell; when the edges size is too big, again, a cell may be crossed in its four edges, this is a situation likely to come out near the function's extreme points and is one of the problems that will be analyzed below). What this second step does is to detect an iso - contour starting point (by locating any crossed edge resulting from step 1, and giving preference to boundary edges to properly locate possible incomplete iso - contour lines in the region defined by the grid, as these have a well defined starting point at the boundary), and follow the contour line through adjacent cells to the one containing the starting point. This operation of following up the line is an iterative one which is done as follows:

- 1.- Having located the edge containing the starting point, set it up as "current edge".

- 2.- Delete any record regarding the "current edge" as being crossed to avoid coming back to it thus building up endless loops.

- 3.- Locate an iso - contour point in the "current edge"

by means of some interpolation procedure between edge extreme points; save it as next contour point.

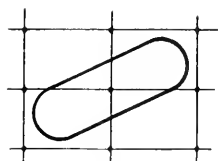
4.- If a "current cell" exists, go to point 5; otherwise decide on one grid cell to which the "current edge" belongs (in general, there will be a choice among two, unless that edge is at the boundary; either one will do); call it the "current cell".

5.- Of the four cells adjacent to the "current cell", pick up the one having the "current edge" in common with it (notice that this cell will have either none or one crossed edge, as crossing information about the common edge was previously deleted); call it the new "current cell".

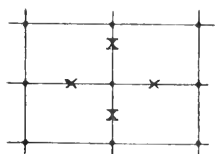
6.- Check the edges in the "current cell" for crossing information; if a crossed edge is detected, call it the new "current edge" and go to point 2. When no crossed edge is detected, the iso - contour is complete. Record this fact and go to point 1 to look for other possible iso - contour branches.

Such an iterative process will be better understood by following the sequence of sketches in Fig. 2.2, where (I) shows the actual iso - contour being built and successive ones refer to procedure steps (a x stands for a crossed edge, a dot for an iso - contour point in the resulting set of points; the shaded cell is the "current cell", the darkest edge the "current edge").

At this point, we are in a better position to discuss the kind of problems that are likely to arise in the neigh



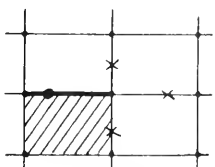
(I)



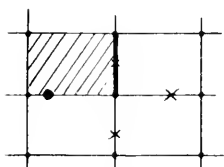
Result of step 1.



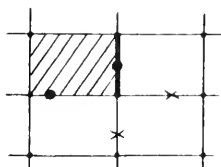
(1)



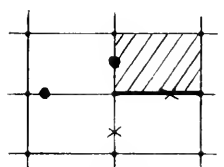
(2,3,4)



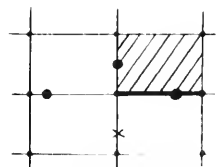
(5,6)



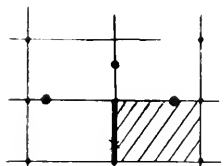
(2,3,4)



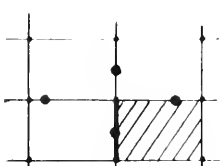
(5,6)



(2,3,4)



(5,6)



(2,3,4)

Fig. 2.2

borhood of function's extreme points. We will discuss three of them.

As has been emphasized repeatedly, grid cell size is critical; near extreme points this is even more true.

It may happen that, the grid size being too large, an small iso - contour near a function extreme point will be roughly contained in one grid cell, as depicted in Fig. 2.3.

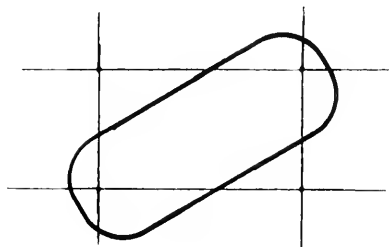


Fig. 2.3

Such a situation violates the assumption that either none or at most two edges were crossed on any given cell. There is no way of deciding in which way the four points in the cell should be joined to generate the iso - contour curve. There are three possible ways of joining them as shown in Figure 2.4. The way that eventually will be used by the procedure

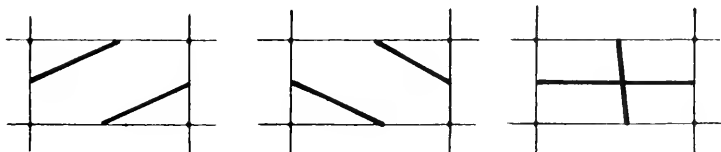


Fig. 2.4

described above depends on a variety of factors, such as on whether it picks one of the cell four crossed edges as star-

ting edge or not and the implicit sequence on which cell edges are checked for crossing information upon having it as current cell.

No reasonable way of solving this problem without changing the cell size (so that the four points will no longer belong to the same cell) was found, unless assumptions are made about topological properties of the function being analyzed. In our case, it is not safe to make any of these assumptions - at least in principle. In addition, given that one of the goals of our study is to determine extreme points, a finer grid is needed in any case to obtain precise estimates. Hence we decided to apply a finer grid whenever this problem appeared. Curiously, this is another situation which points out the weakness of software procedures in dealing with pattern recognition; most of the early approaches to such problem were very similar ;i.e., superimposing a grid to the object under inspection and being unable to infer details from the general pattern. In our case, it is obvious that a human being, with the help of nearby iso - contours would have no problem in joining the four points in that cell in the correct way. The procedure doesn't generate such information, as it constructs one iso - contour at a time, but there is no apparent way to make ^(it) available to the routine to solve the problem unless very inefficient **procedures**, such as slope matching are used.

A second problem, particularly relevant when the function under study is very flat near the extreme points (and this was our case, as it will be seen later), is to get

a set of adjacent grid points all with the same function value, in turn equal to the level of the iso - contour being constructed. The situation is illustrated in Figure 2.5. Assume the values besides grid points are the function values at them, and that we are building the iso - contour of level 5. What

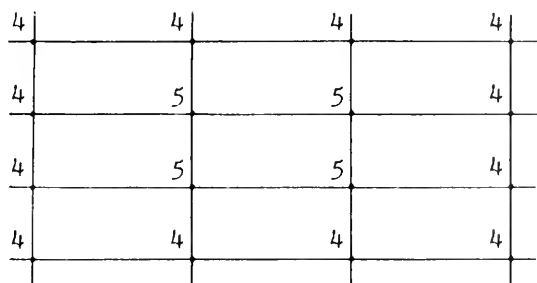


Fig. 2.5

seems logical to do is to record the edges with the maximum value at their extreme points as being crossed, since they themselves are so closed to the actual iso - contour. Test (a) above would indeed mark those edges as being crossed, but it would as well mark all the other edges with an extreme point of level 5. Then, what may happen (depending upon the same kind of factors outlined for the previous problem), is that we may end up with an iso - contour plot as ridiculous as the one depicted in Figure 2.6, consisting of four independent points considered as iso - contour disjoint branches.

Reference (2) proposed a way of avoiding this problem by disturbing the function values slightly in the proper direction so that only the outer edges would be recorded

as crossed (notice that when those points are at a relative minimum the disturbance has to be in the opposite direction).

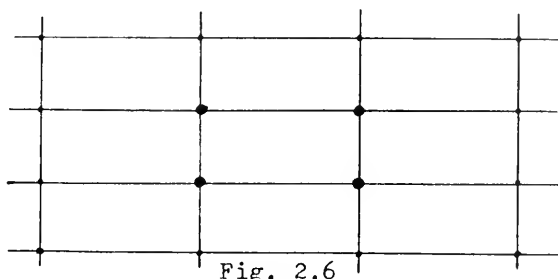


Fig. 2.6

This preserves the nature of test (a), but search for such grid nodes has to be undertaken as a pre-step for the method in use. The way we solved it is by complementing (a) in such a way that edges with extreme values both equal to the iso - contour level are not recorded; only those with a non zero difference at one of their extreme points are so. It is equivalent to the other one and it does not require any pre-step.

Finally, the problem depicted in Figure 2.1 is one which is a direct consequence of the cells being too large. It can not be solved by any other procedure than using a finer grid, since there is no way of detecting two crossing points in an edge with information about edge extreme points (and even if it were, the interpolation process would end up with the same points along the edges generating two exactly equal branches for the iso - contour).

3.- Asymptotic expansion behavior and maximum likelihood estimation.

We had at our disposal two sets of sample data to study the behavior of the likelihood function generated by them: the first is a set of 24 North Sea pools and the second is a set of 54 pools from a play in the Western Canadian Sedimentary Basin.

Before going any further, it is important to say that the chronological order in reservoir discoveries is of central importance, as the sampling process is assumed to be without replacement and proportional to -remaining- size.

Some of the results obtained using the expression of the likelihood function on those samples are summarized in tables 1 and 2. In them, the values of μ and σ^2 which maximized the likelihood function for different values of N and sample size n are shown (When the sample size is less than the actual value stated above, it has to be interpreted that the sequence of first sample values was used as data, thus preserving chronological order). The estimates for N = ∞ were computed analytically, that is,

$$\sigma_{\infty}^2 = \frac{1}{n} \sum_{j=1}^n (\log Y_j - \log g)^2, \quad ,$$

$$\mu_{\infty} = \log g - \sigma_{\infty}^2, \quad \text{where}$$

$$\log g = \frac{1}{n} \sum_{j=1}^n \log Y_j, \quad ,$$

and Y_j stands for the j^{th} observation.

TABLE 1
NORTH SEA DATA

n	N	μ_0	σ_0^2	Likelihood value	Figure #
24	200	2.47	1.45	-128.279	1
24	300	2.45	1.6	-128.124	2
24	303	2.422	1.569	-128.12	3
24	400	2.3	1.7	-128.179	4
24	500	2.32	1.6	-128.269	5
24	600	2.35	1.6	-128.369	6
24	750	2.5	1.3	-128.463	7
24	999	2.6	1.2	-128.314	8
24	2000	2.6	1.2	-128.894	9,10
24	5000	2.6	1.2	-127.551	11,12
24	10000	2.6	1.21	-127.425	13
24	20000	2.6	1.21	-127.362	14,15
24	10^5	2.59	1.22	-127.31	16,17
24	∞	2.50286	1.27988		
20	300	1.95	1.95	-103.827	18
20	400	1.75	2.15	-104.01	19
20	∞	2.334	1.31567		
12	200	0.	5.2	- 59.7028	20
12	300	0.	5.	- 60.0069	21
12	∞	2.1806	1.48137		

TABLE 2

WESTERN CANADIAN SEDIMENTARY BASIN DATA

n	N	μ_0	σ^2	Likelihood value	Figure #
52	500	2.2	2.3	-314.024	22
52	1000	1.85	2.5	-310.084	23
52	1500	1.7	2.6	-308.58	24
52	2000	1.52	2.7	-307.813	25
52		1.2648	2.72134		
40	1000	1.85	2.8	-245.863	26
40	1500	1.5	3.	-244.921	27
40	2000	1.3	3.2	-244.442	28
40	3000	1.	3.5	-244.012	29
40		1.40947	2.95797		
30	1000	0.8	3.9	-183.915	30
30	1500	0.1	4.7	-183.342	31
30		1.26892	2.95797		
20	1000	-4.	12.	-126.765	32
20	1500	0.	5.7	-129.761	33
20		2.02979	2.74941		

The column labeled "Figure #" indicates the iso - contour plot corresponding to each row. In these plots, the axis labeled "SIGMA" actually corresponds to σ^2 values.

We will center the discussion in the values shown in Table 1, as we have a wider range for them and the operating characteristics of the asymptotic expansion for the likelihood function which we detected seem to apply to the values on Table 2 as well.

The main feature to be pointed out is the degree of instability which shows up for μ and σ^2 estimates as N stays constant and the sample size is changed. More concretely, as the sample size n increases, ML estimates (keeping N fixed) for μ and σ^2 vary systematically: the estimate for μ increases while that for σ^2 goes down.

Furthermore, as n takes values below a certain level, such changes in μ and σ^2 estimates become more apparent, taking values very different from the ones obtained for larger values of n. This is the case for n = 12 in Table 1 (even for n = 15 with those data such a behavior begins to be very apparent) and for n = 30, 20 in Table 2.

There are, as we can see, two two main reasons for such instability.

The first is probably a consequence of sample data variability, which is relatively more important when the sample size is small. The characteristics of the samples we had -perhaps due to the fact that oil producing areas are so recognized when important discoveries take place, so that some early observations tend to be large- was such that

when used in the likelihood function produced a bimodal behavior which was hard to detect as the second relative maximum builds up in a region where the μ and σ^2 values are unreasonable as estimates of the underlying density.

Moreover, the second relative maximum tended to be more important as N increased for fixed n . As may be seen in the plots corresponding to $N = 5000, 20000, 100000$ when $n = 24$ for the first sample (Figures 11, 12, 14, 15, 16, 17), it reaches a point where becomes more important than the first one. When this happens, those values out of the reasonable ranges for μ and σ^2 mentioned above begin to appear.

For very small samples (e.g., 12 in Table 1, 20 or 30 in Table 2), it seems that the second relative maximum has taken over completely, and it is the only one obtained in those cases.

Before discussing the other reason for instability, more intrinsic to the nature of the problem and less related to data variability, it is perhaps worthwhile to make a methodological point for using the iso - contour routine or a similar tool in numerical analyses like the one described here. We would recommend a wide range on the function arguments to take the first pictures, forgetting about "reasonable" values around which the function is expected to behave in a certain way (e.g., to have a maximum). Having an overall idea of function behavior we can then concentrate on certain areas. If this step is omitted, the first results tend to bias the analysis to pursue a certain region without paying any attention to others which may contribute heavily to explain peculiar be-

avior. In our environment, doing so implied a loss of detail, as the plots we generated were fixed in size, but even in this case it proved to be in the right direction.

The second cause of estimate instability was due to the extreme flatness of the function near its maximum(s). As it may be seen in the enclosed plots, this circumstance did show up rather strongly. Thus, small variations in the functional form due for example to small changes in sample data are likely to produce relatively important changes in the location of the function's extreme point(s).

The changes observed in the location of maximums were not, however, completely arbitrary. What happened was that the iso - contours we generated were roughly elliptic, and that the main axes of such ellipses tended to conserve their orientation in the $\mu - \sigma^2$ plane. More concretely, the main axis tended to conserve its location, so that different estimates tended to lie in the line defined by it. This may be seen as a common characteristic of all the plots we include.

As a function of N with n fixed, the iso - contour plots show that the function gets tighter along the small ellipses' axis, but being still very flat along the other one. This suggested that the underlying colinearity between μ and σ^2 estimates in the limit as $N \rightarrow \infty$ (see μ_∞ and σ_∞^2 equations above) could be somewhat preserved for finite values of N . Visual inspection of the plots pointed out a line equation close to

$$\mu = 3.775 - 0.8 \sigma^2$$

for the first sample as main ellipses' axis. In the limit, for

$n = 24$, the equation is

$$\mu_{\infty} = 3.78274 - \sigma_{\infty}^2,$$

not very far off the one visually fitted. This would explain the overall trend of μ and σ^2 estimates as the sample size n changes which we described at the beginning of this section ; it may help to foresee their behavior for different sample sizes.

We were also interested in obtaining an estimate for N , the finite population size. Our results in this regard are not very conclusive. As shown in Table 1, consulting the column labelled "Likelihood values", it appeared as if for $N = 303$ we obtained an absolute maximum. However, as we had other results for increasing values of N , the function value increased above the quantity -128.12 corresponding to $N = 303$, and, what is worse, the overall maximum (up to -116.511 for $N = 100,000$ in Figure 16) is attained at a point with unreasonable values for μ and σ^2 . Thus, at least for samples as small as the ones we had, ML estimators for N appeared to be quite instable, and other estimation methods should perhaps be tried to obtain a value for the population size. It has to be seen, however, how the asymptotic expansion works with regard to N for samples with a larger n . It may well be the case that our sample sizes were too close to values for which the second function's relative maximum begins to develop, thus jeopardizing the function usefulness in the relevant range for μ and σ^2 .

To summarize, the two reasons for estimates' instability discussed above have pointed out how sample size

influences the estimates we obtain with the analyzed likelihood function expansion. Its pathological behavior for small samples seems to suggest that having more numerous samples would help to obtain better estimates, not only because they provide more information, but also because the likelihood function begins to behave better when n increases.

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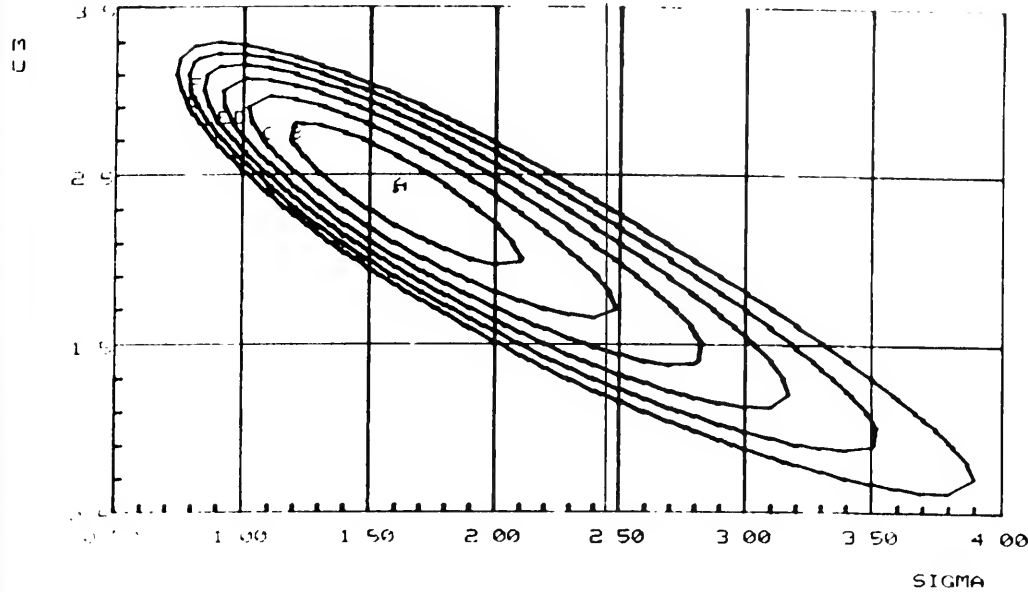


Figure 2.

ANS08303/FIRST 3 CONTOUR VALUES ARE (THE REST AS IN FILE LEVEL305):
 A=-128.12,B=-128.2,C=-128.3.

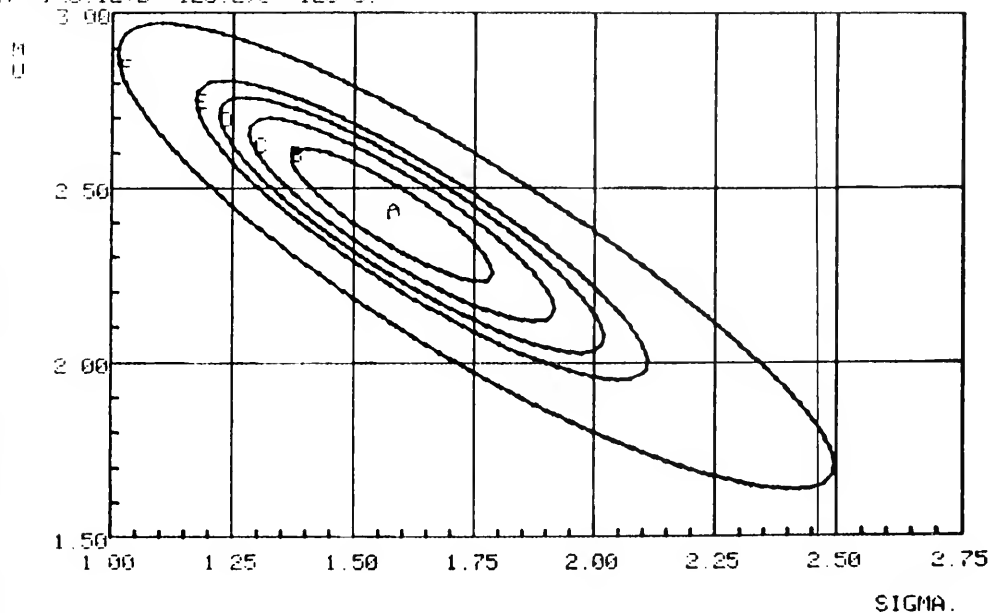


Figure 3.

INSOB400/ THE FIRST 3 CONTOUR VALUES ARE (THE BEST AS IN FILE LEVEL400)
 A=-128.179 B=-123.510 C=-129

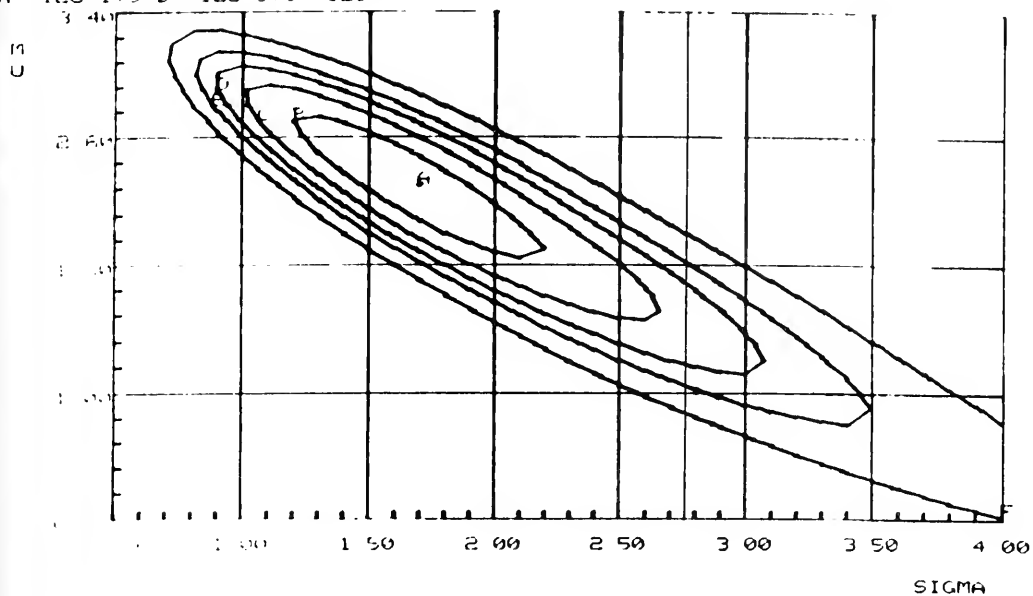


Figure 4.

PH08500>FIRST 3 CONTOUR VALUES ARE THE BEST HS IN FILE NI24500>
 AH=108.269, B=-138.5, C=-129

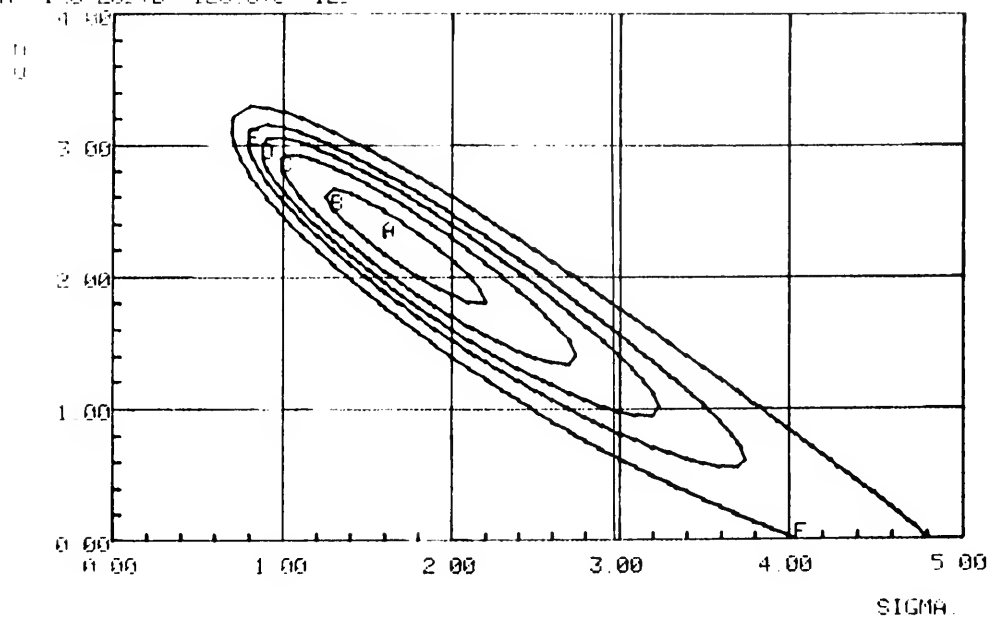


Figure 5.

INSOB600/ THE FIRST 3 CONTOUR VALUES ARE (THE BEST RS IN FILE LEUPLAND)
 A=-128 B=0 C=-128 S.C=-120

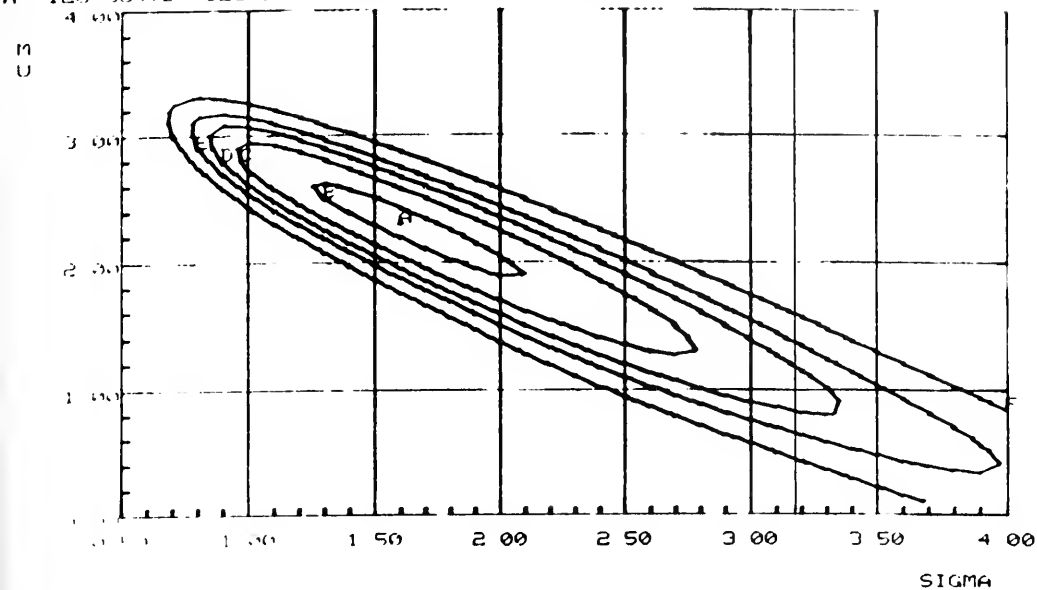


Figure 6.

DS08750: FIRST 3 CONTOUR VALUES ARE (THE REST AS IN FILE LEU24750)
H=-128.463, B=-128.75, C=-129.

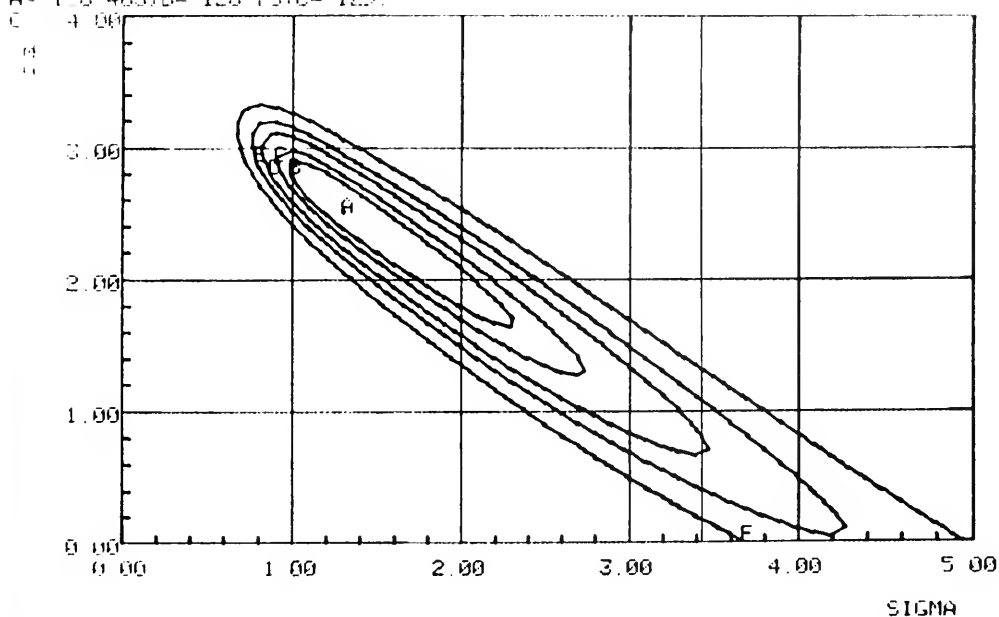


Figure 7.

PROB999 THE FIRST 3 CONTOUR VALUES ARE THE BEST NS IN FILE 11041009
 G=-120.014 E=-120.510 *-120

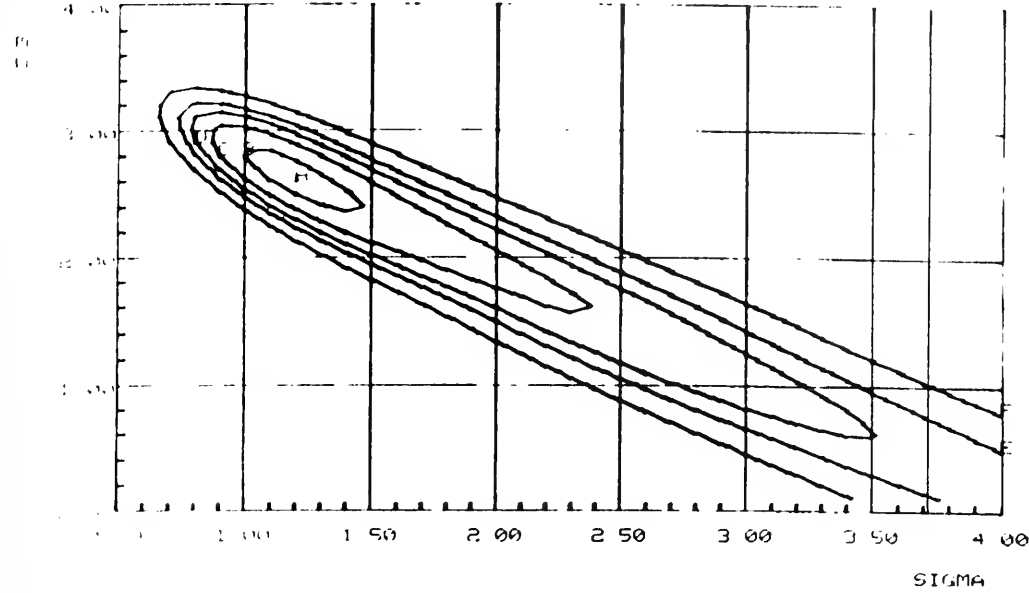


Figure 8.

HS2000: FIRST 3 CONTOUR VALUES ARE THE REST AS IN FILE LEN2000).
 A=-127.694 B=-128.5 C=-129.

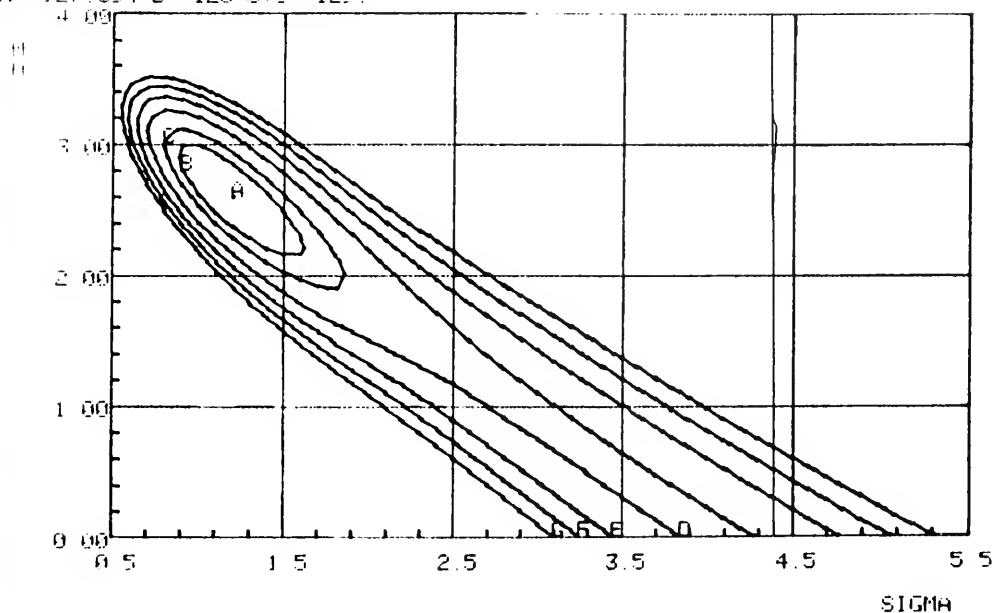


Figure 9.

>NE2000>FIRST 3 CONTOUR VALUES ARE THE BEST AS IN FILE LEU2000>
 H=-17.894, B=-128, S, L=-129

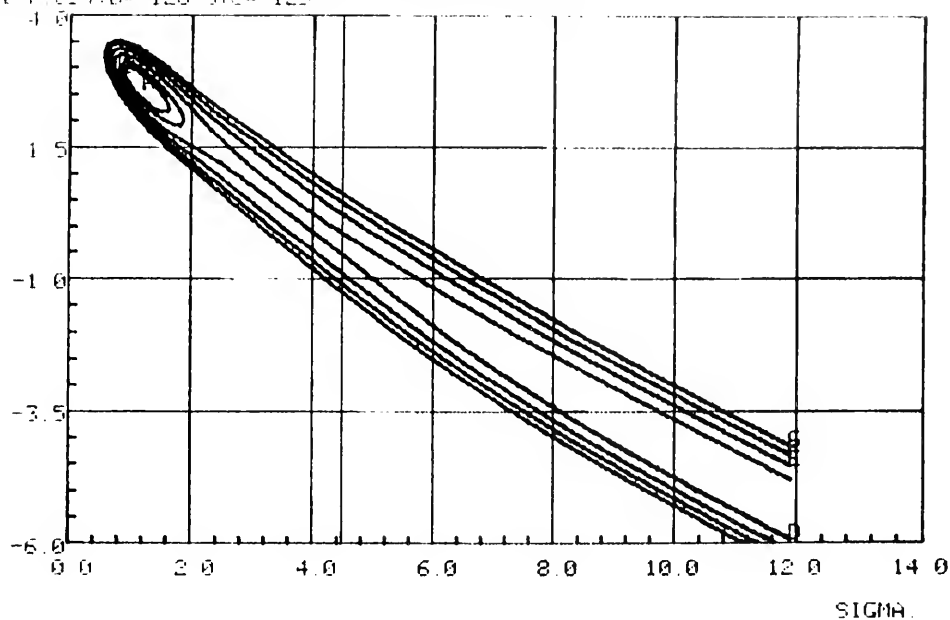


Figure 10.

LEVS0000/ THE FIRST 3 CONTOUR VALUES ARE THE REST AS IN FILE LEV5000 >
H=-127.551,B=-128.0,C=-128.5

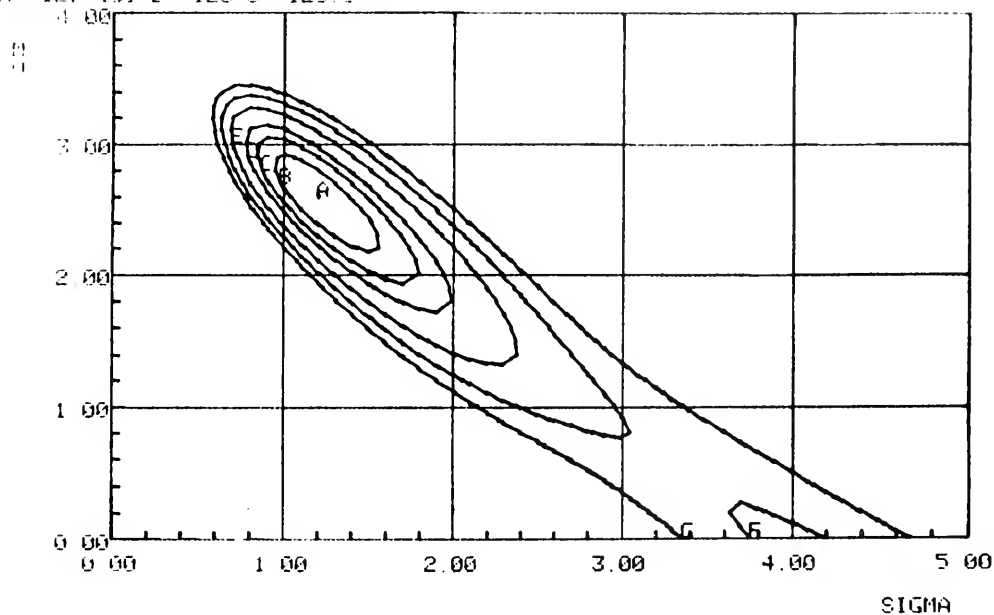


Figure 11.

LEU5000/FIRST 3 CONTOUR VALUES ARE(THE REST AS IN FILE LEU5000):
 A=-135.357,B=-126,C=-127.
 CLOUDS COMPARISON:

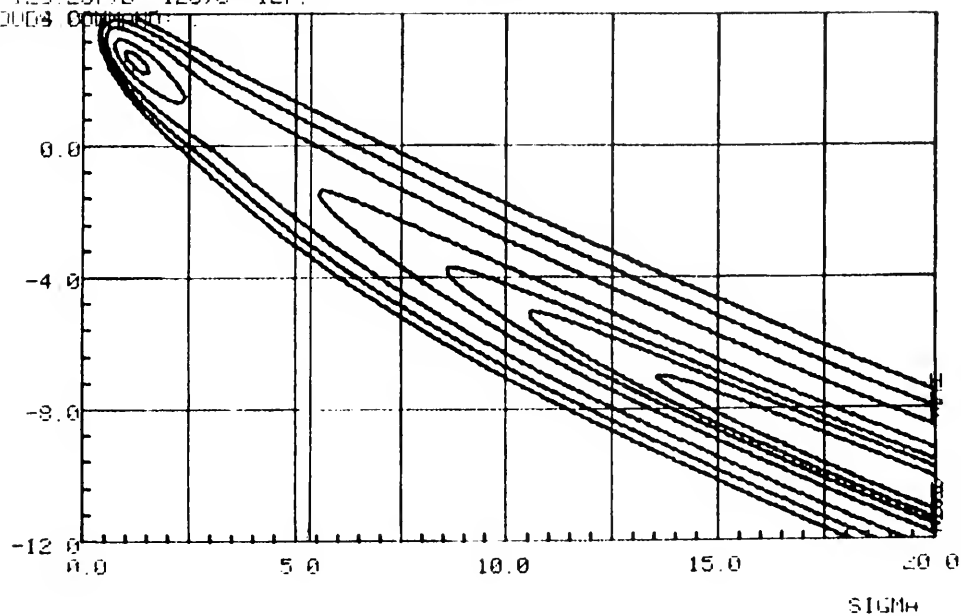


Figure 12.

NS10000 THE FIRST 3 CONTOUR VALUES ARE THE REST AS IN FILE LE100000
 A=-127.425 B=-127.2 C=-122.5

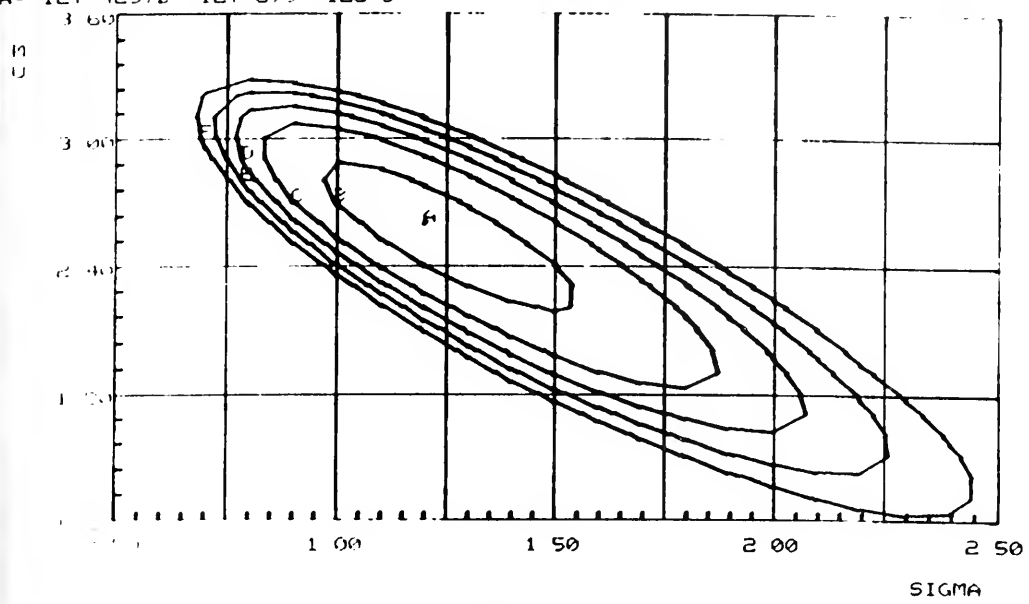


Figure 13.

LE200000 THE FIRST 3 CONTOUR VALUES ARE THE BEST AS IN FILE LE200000
 A=-127.362/B=-127.8/C=-128.5

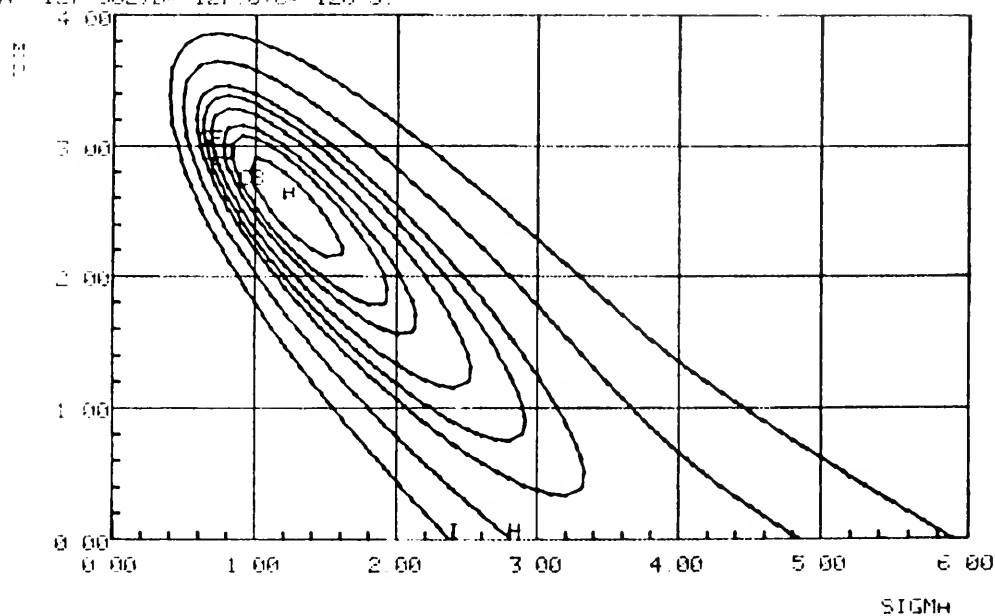


Figure 14.

LE20000 (FIRST 3 CONTOUR VALUES ARE THE REST ARE IN FILE LE20000)
 A=-119.691,B=-121.0,C=-125

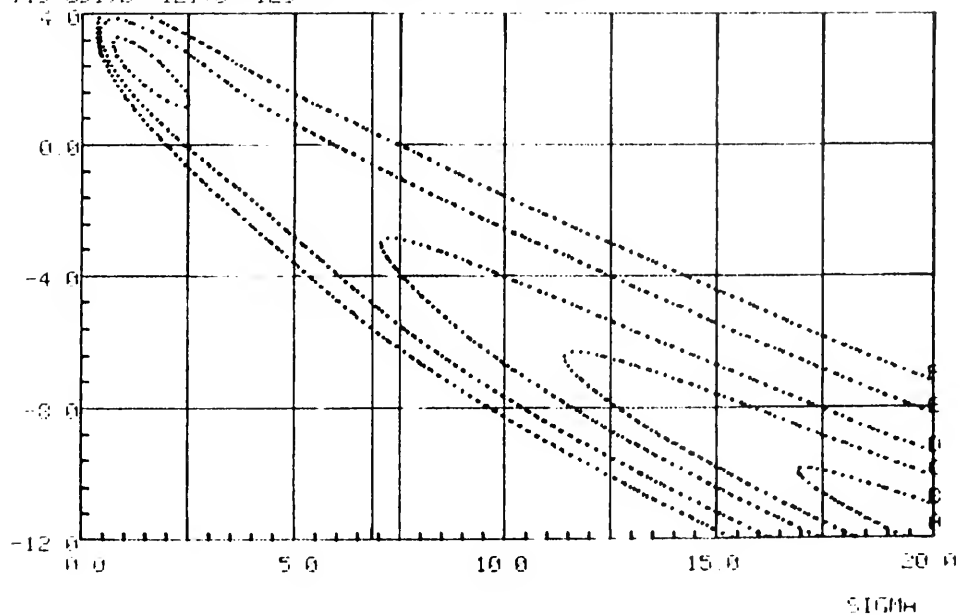


Figure 15.

LE100000 FIRST 3 CONTOUR VALUES ARE (THE REST AS IN FILE LE100000)
RA=-127.31, B=-127.8, L=-129.5.

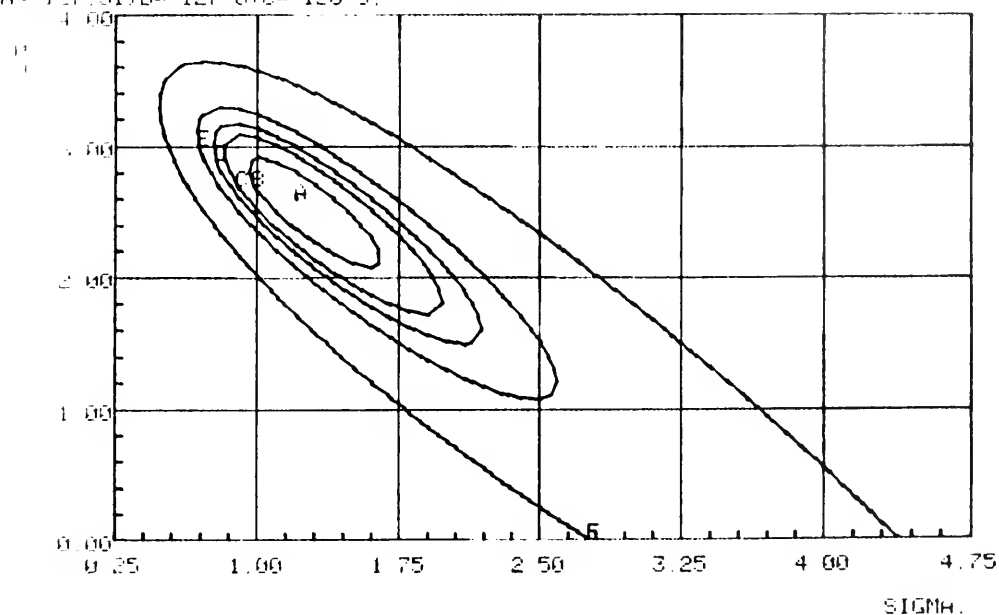


Figure 16.

1110000 (FIRST 3 CONTOUR VALUES ARE THE REST AS IN FILE LE100000)
 H=-115, G11, B=-118, C=-122.

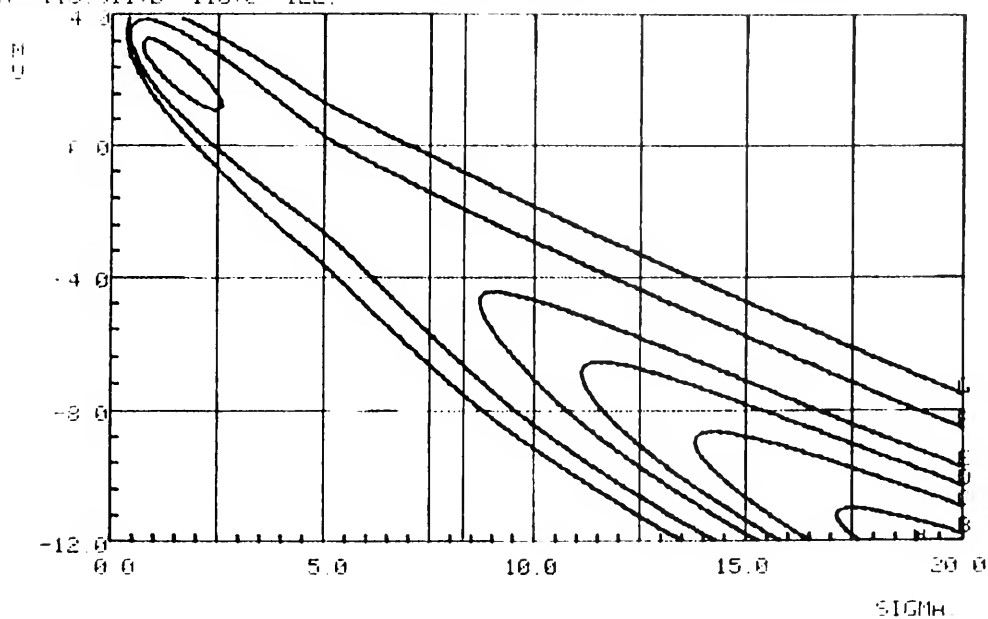


Figure 17.

ANS20300/FIRST 3 CONTOUR VALUES ARE (THE REST AS IN FILE LEU20300):
 A=-103.827, B=-103.9, C=-104

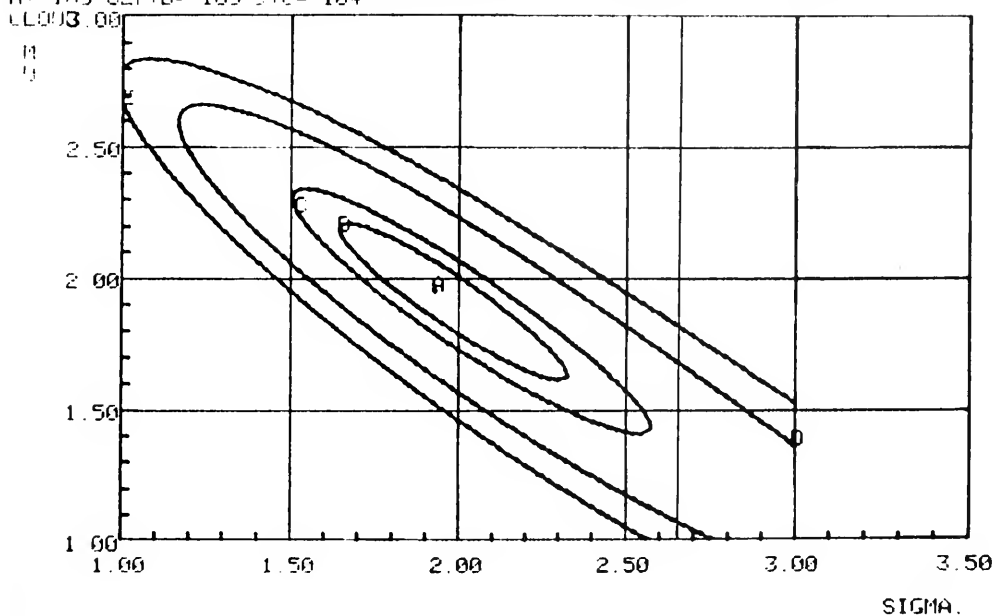


Figure 18.

NS20400/FIRST 3 CONTOUR VALUES ARE THE BEST AS IN FILE LEU20400
 1.0×10^{-4} 0.1×10^{-4} 0.01×10^{-4}

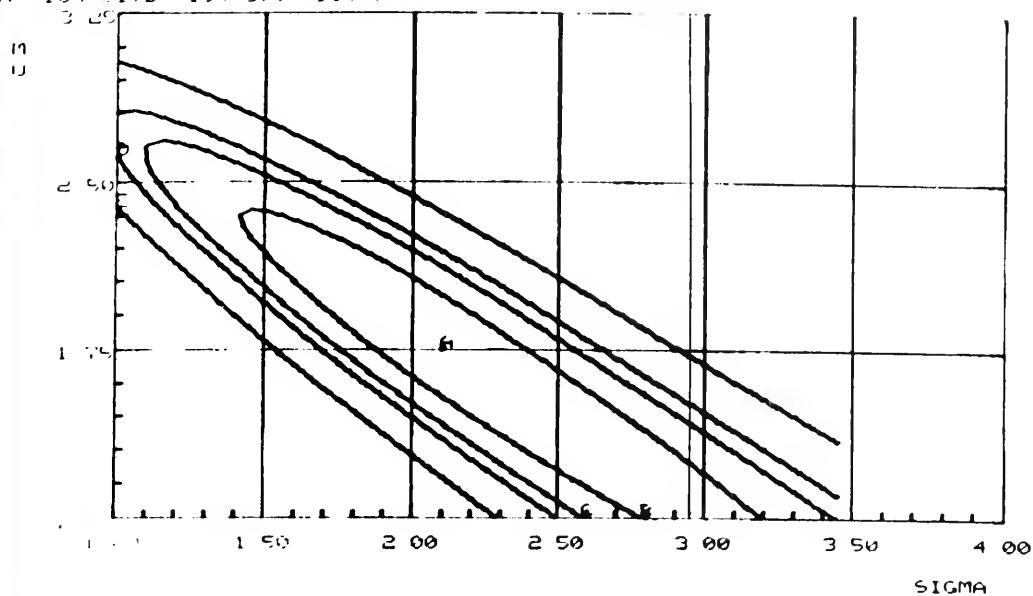


Figure 19.

(HIS12200/FIRST 3 CONTOUR VALUES ARE (THE REST AS IN FILE LEV12.200))
 H=-59.7020,B=-60.00-00

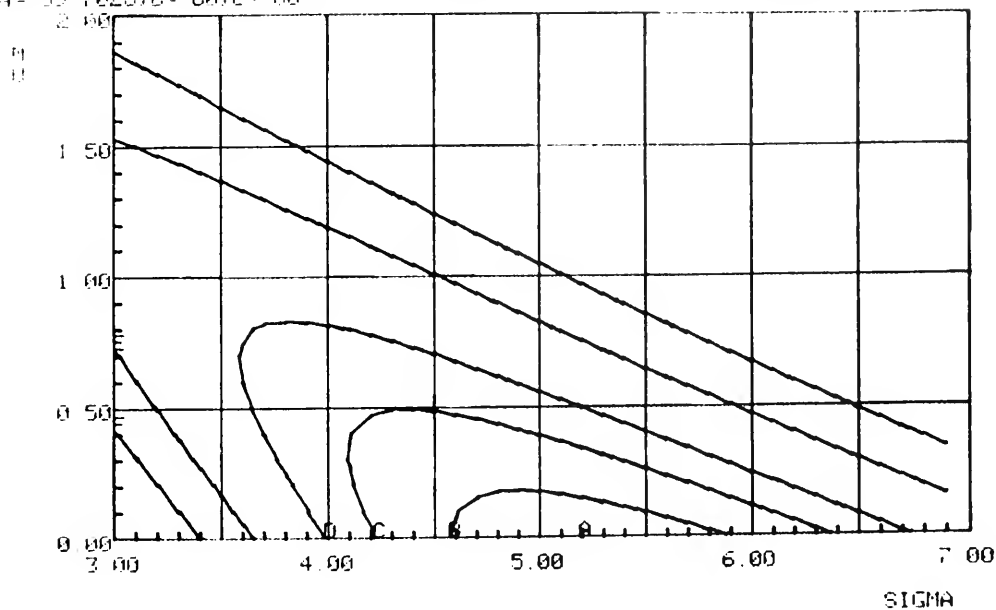


Figure 20.

NS12300/FIRST 3 CONTOUR VALUES ARE THE REST AS IN FILE LE112.300)
 A=-60.0069, B=-60.5, L=-61.

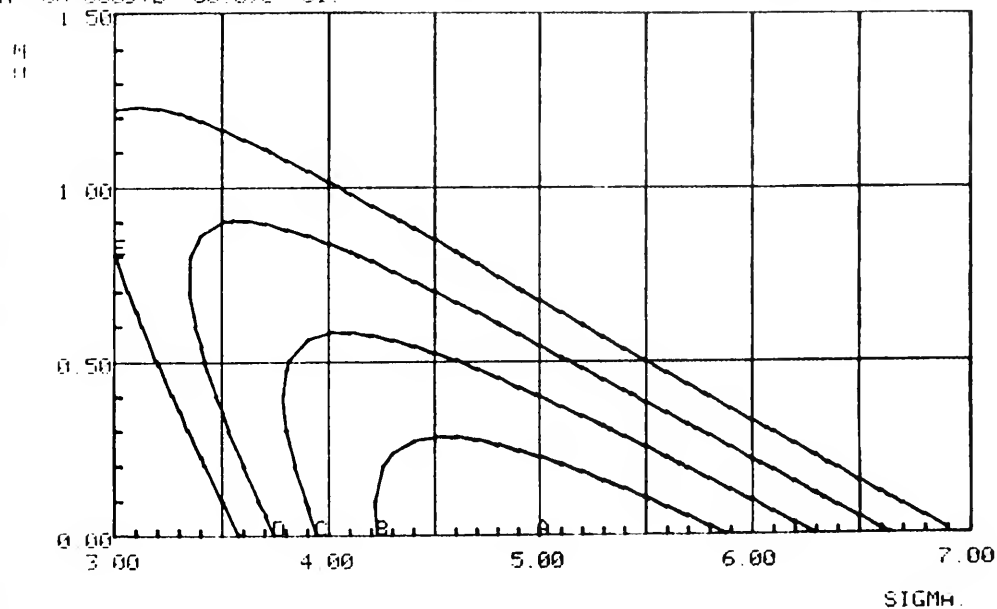


Figure 21.

.D52500/ THE FIRST 3 CONTOUR VALUES ARE (THE REST AS IN FILE N102500)
 A=-314.023.B=-314.5.C=-315

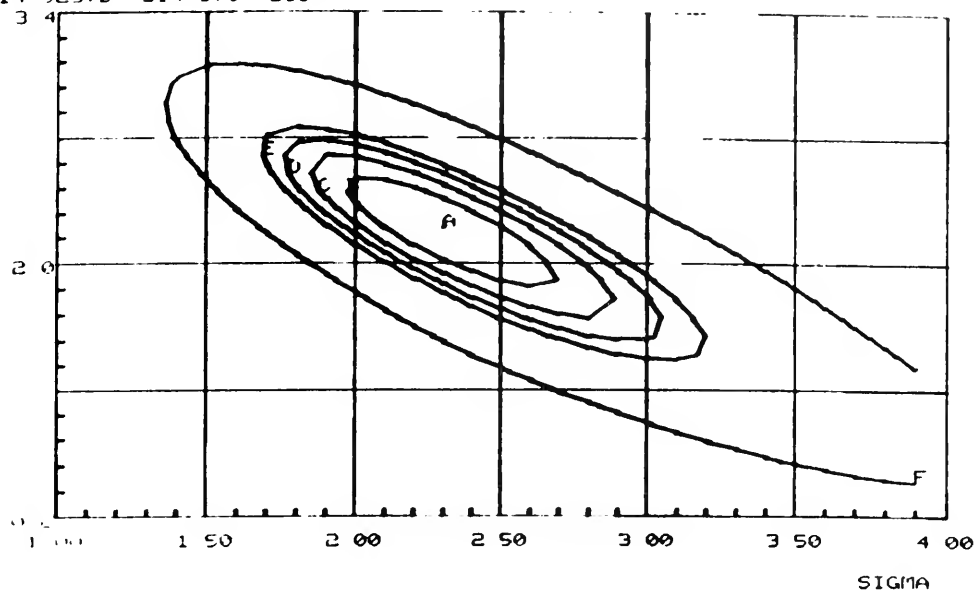


Figure 22.

0521000 THE FIRST 3 CONTOUR VALUES ARE (THE BEST AG IN FILE H1521000)
 H=-310.034, B=-310.1, C=-310.2

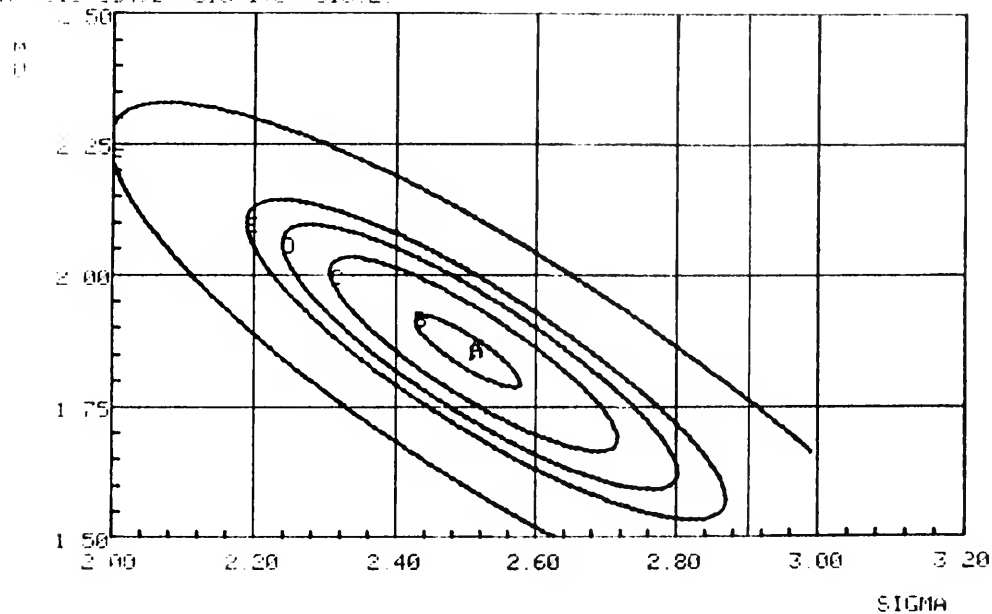


Figure 23.

DIS 1500 THE FIRST 3 CONTOUR VALUES ARE THE BEST AS IN FILE HT11500
 (M=50.0E+009.0E+009.5)

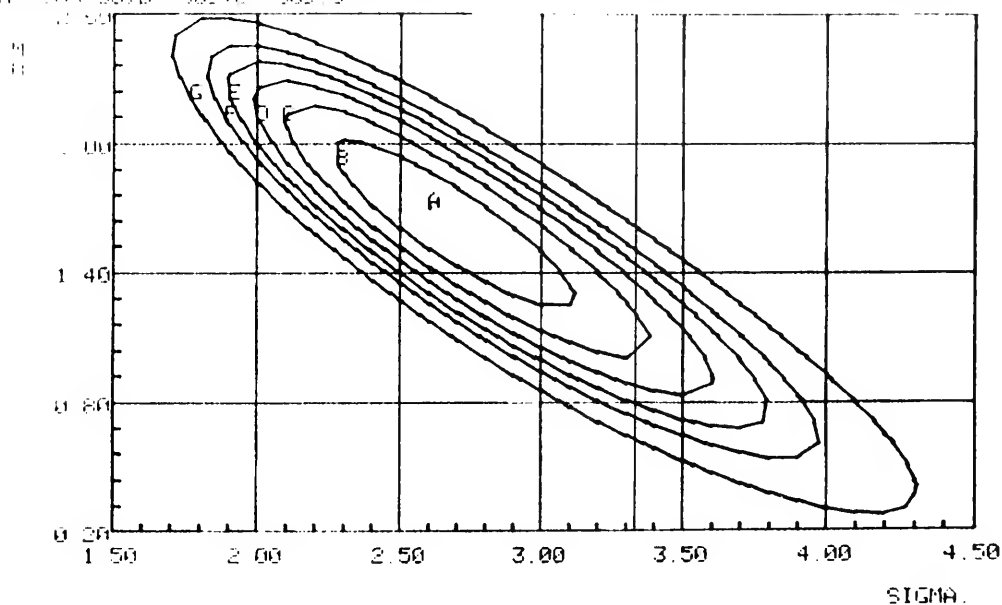


Figure 24.

405,0000 FIRST 3 CONTOUR VALUES ARE THE BEST AS IN FILE H1522000.
 A=-307.813; B=-308.0; C=-309.5

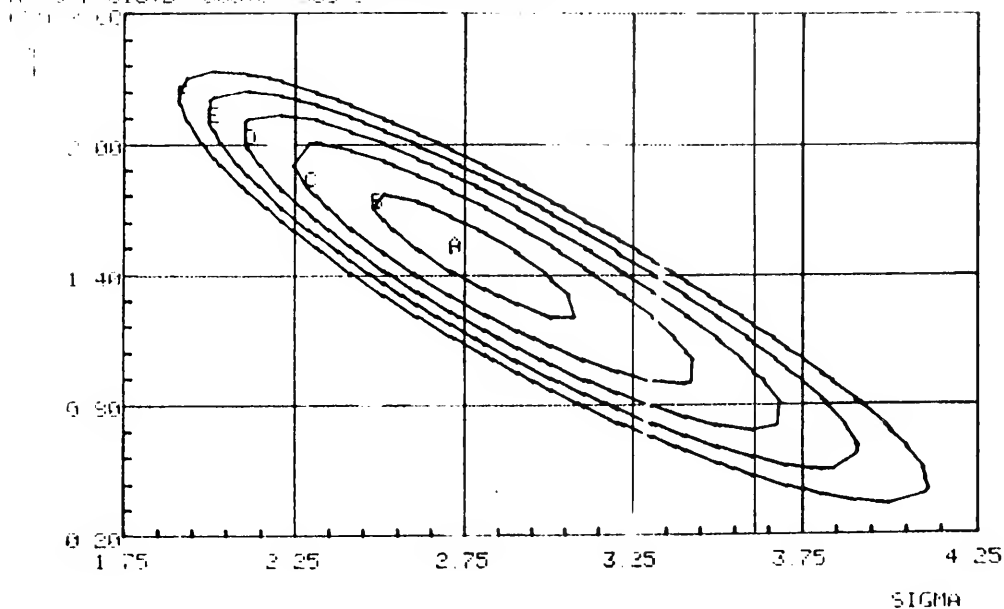


Figure 25.

0401000:FIRST 3 CONTOUR VALUES ARE THE BEST AS IN FILE 01401000):
H=-245.863,B=-246.0,C=-246.5

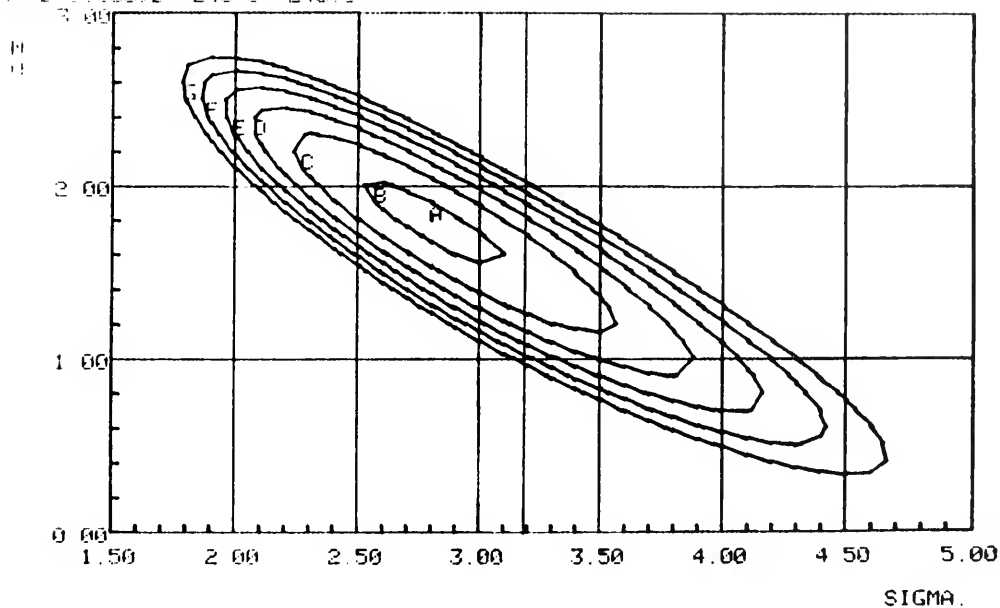


Figure 26.

0001500: FIRST 3 CONTOUR VALUES ARE THE BEST AS IN FILE HI401500 +
 0001500: 14 921 10=-245 20=-245 5

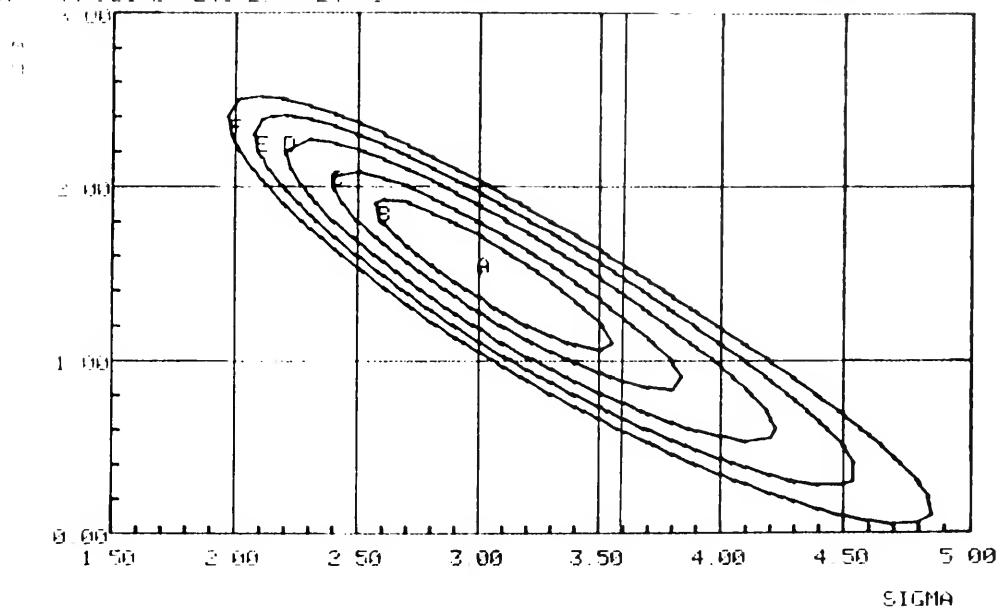


Figure 27.

(H402000) THE FIRST 3 CONTOUR VALUES ARE THE BEST AS IN FILE H402000
 H=-244.442, B=-244.6, C=-244.0.

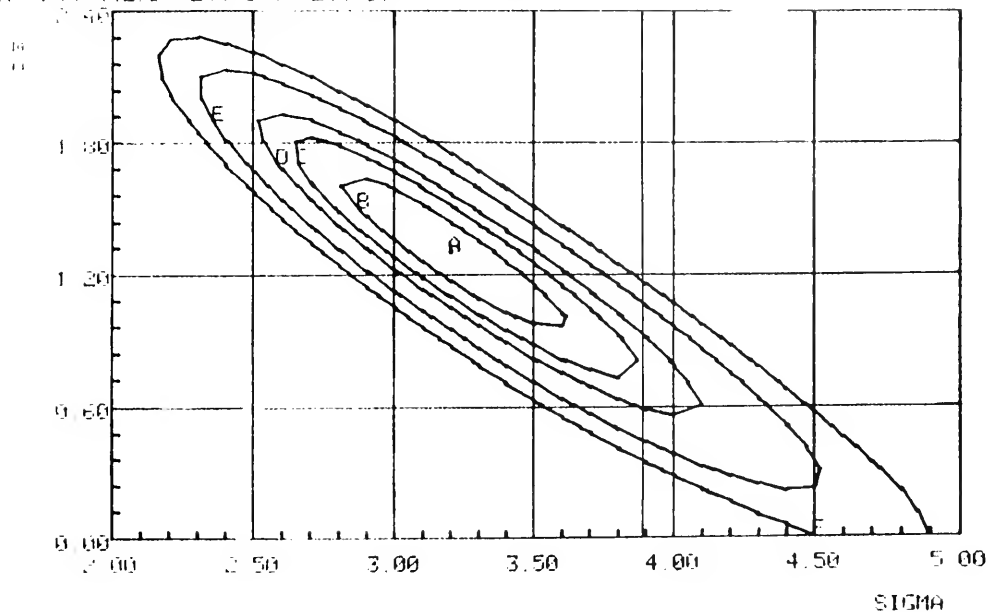


Figure 28.

NI403000 (THE FIRST 3 CONTOUR VALUES ARE THE REST AS IN FILE NI403000)
 ME=-244.01C=-244.24E=-244.50

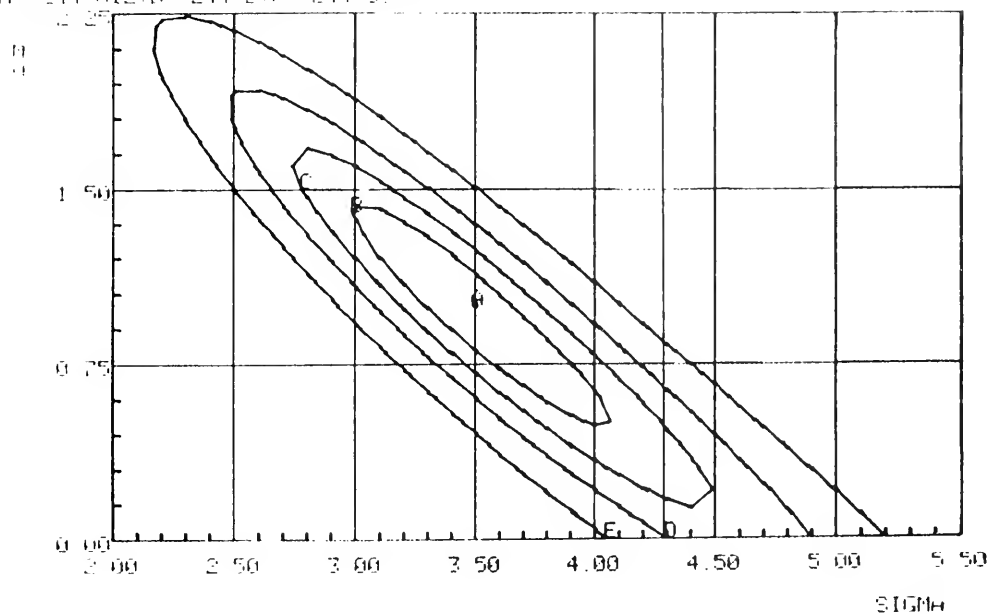


Figure 29.

0301000 FIRST 3 CONTOUR VALUES ARE THE REST AS IN FILE HI301000
 RE=187.915, B=-184.3, L=-184.5.

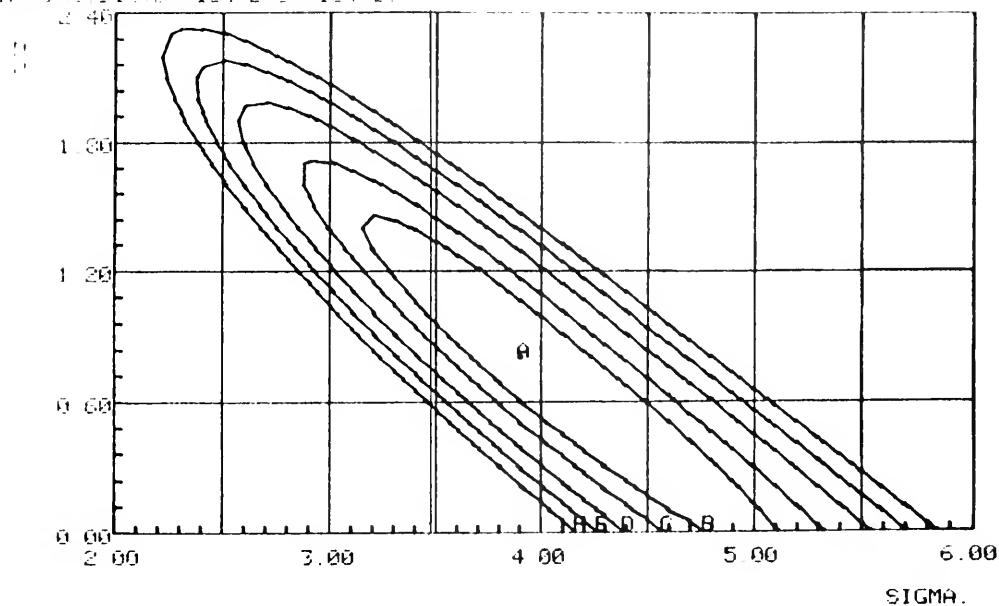
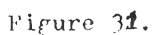


Figure 30.

8



0000000/ THE FIRST 3 CONTOUR VALUES ARE (THE REST AS IN FILE H120MILA):
 HT=126.765, B=-127.0, C=-127.5

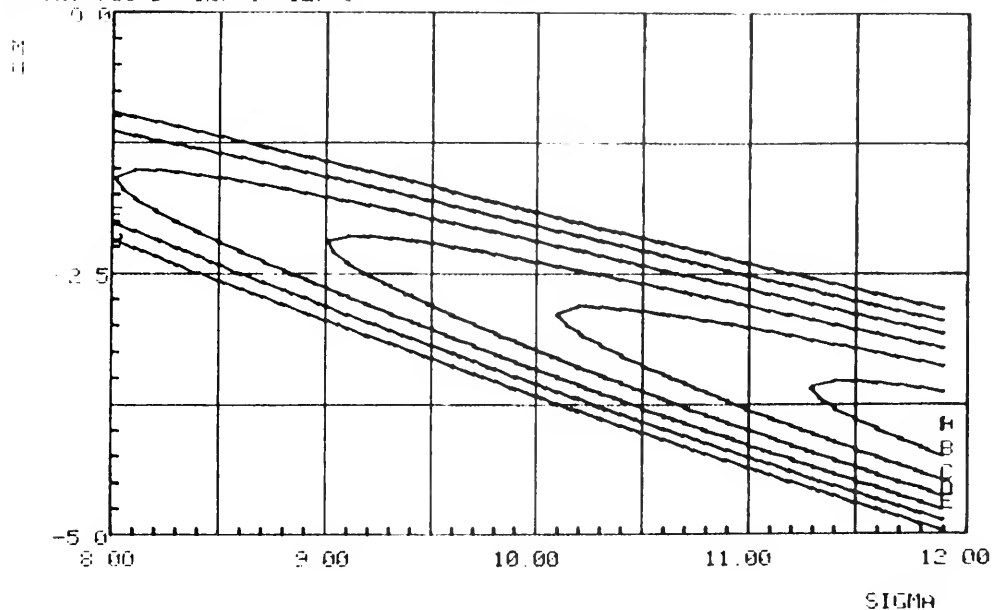


Figure 32.

000500 FIRST 3 CONTOUR VALUES ARE THE REST AS IN FILE N1201500
 H=-130.761,B=-130.0,-130.5

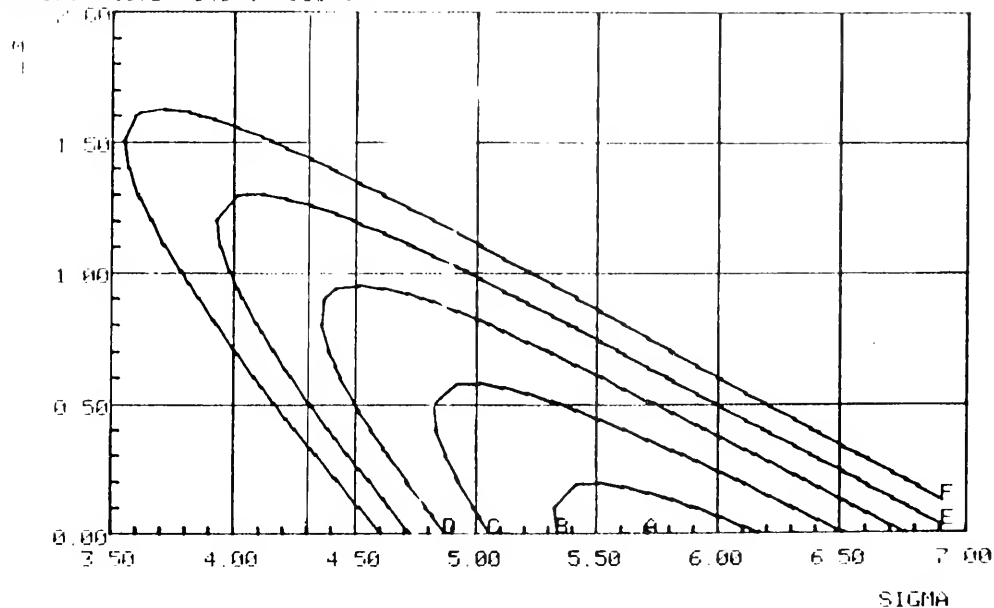


Figure 33.

APPENDIX AIso-contour plotting routine description and program listings.

The iso-contour generating routine consists of a set of FORTRAN and 370 Assembler subroutines which form a package. CTOUR is the name of the main routine in this package, and its purpose may be generally described as follows:

Given a tabulated function of two independent variables $f(x,y)$, it determines sets of points belonging to the locus defined by the equation $f(x,y) = C$, where C is a given constant. (Subsequent calls to CTOUR with different C values allow to obtain a set of such loci).

1.- Procedure.

CTOUR employs as main procedure the one described in (1). Roughly, it works from a table of function values at points of a rectangular grid defined by two auxiliar input arrays (which permit the use of variable increments in x and y along their respective axes). Given a value for C in $f(x,y)=C$, CTOUR begins by marking all the mesh edges crossed by the locus defined by the above equation (an auxiliar bit memory is used to keep track of edges); then it computes the coordinates of the crossing points by linear interpolation between function values at adjacent grid points. Points belonging to the locus are returned as result, ordered in such a way that it is possible to actually draw the contour by just joining them

(1) Automatic Contour Map. G. Cottafava and G. Le Moli, Comm. of the ACM, July, 1969.

as they are found in the result array. This result array may be logically divided in subtours, when the locus $f(x,y) = C$ is such that it possesses disjoint branches in the domain defined by the grid.

2.- Input and output variables description.

It is assumed that CTOUT will be called from a FORTRAN program. Its calling sequence is as follows:

```
CALL CTOUT(A,NR,NC,XS,YS,ALEVL,IA,PTS,NDPTS,NPTS),
```

where:

- A is a bidimensional array containing the tabulated function values. By convention, X values are constant along A's columns, while Y values are constant along A's rows. For convenience, A must be bordered by two extra rows and two extra columns which do not need to contain function values, but which must be included in the count to determine A's dimension.
- NR is the number of rows in A, as declared in the calling program (including the extra first and last rows).
- NC is the number of columns in A as declared in the calling program, also including the 2 extra columns.
- XS is an unidimensional array with NC elements containing the X values corresponding to different columns of A.
- YS is another unidimensional array with dimension NR containing the Y values corresponding to different rows of A. (2)

(2) In other words, $A(I,J) = f[XS(J),YS(I)]$.

- ALEVL is the function value for which the contour points are to be determined.
- IA is an auxiliary memory taking the form of an unidimensional integer array. It must contain at least $\left\lceil ((NR-1)*NC+(NC-1)*NR)/32 \right\rceil$ elements, and it must be declared accordingly in the calling program.
- PTS is a bidimensional array where the locus' points are to be returned. It is of the form $PTS(NDPTS,2)$, where $PTS(.,1)$ will contain X coordinates and $PTS(.,2)$ Y coordinates for the resulting locus' points.
- NDPTS is the maximum number of points that may be allocated in PTS (i.e., its first dimension as declared in the calling program). If CTOUR reaches this maximum during the process of filling PTS, it calls another subroutine (EXTEND) which is supposed to save the contents of PTS when enough space to do so is available, cancelling the whole process otherwise. An actual program for EXTEND is not included below because it is likely to vary in different environments.
- NPTS is a result integer variable indicating how many points are returned in PTS.

3.- Example.

In what follows we describe how to set up the input variables for a call to CTOUR in a concrete case, along with the output generated by it.

Imagine that we have a function $f(x,y)$ tabulated in the interval $[1 \leq x \leq 3, 2 \leq y \leq 5]$, by increments of 0.5 in both

x and y, and that we want to call CTOUR to obtain the set(s) of points belonging to the locus of level 4.5, that is, belonging to the curve defined by $f(x,y) = 4.5$.

If that is the case, the input data should be organized as depicted in Fig. A.1, and the grid being used would be the one shown in Fig. A.2.

Assume further that the iso-contour $f(x,y) = 4.5$ has the two branches drawn in Fig. A.2. Then, the result variables returned by CTOUR upon call with the Fig. A.1 variables would be of the form depicted in Fig. A.3. Notice that the first point of a contour branch is repeated as its last one to actually close it, unless the branch is incomplete with starting and ending points at the grid boundaries. Different branches are separated by delimiters labelled NA in both the x and y coordinates.

4.- Brief process description, with references to actual programs.

All subroutines are written in FORTRAN, with the exception of MARK, ERASE and SEE which manipulate bits in the auxiliary memory IA and are written in 370 Assembler for improved efficiency.

Fig. A.4 shows a high level flowchart for CTOUR, indicating calls to other subroutines in the package. It is quite straight forward, although some steps need further explanation. Following are succinct discussions of such steps and the involved subroutines.

- "Clear auxiliary memory" just sets all the integer elements in IA to zero.

IA is used as a bit array, each bit corresponding to one

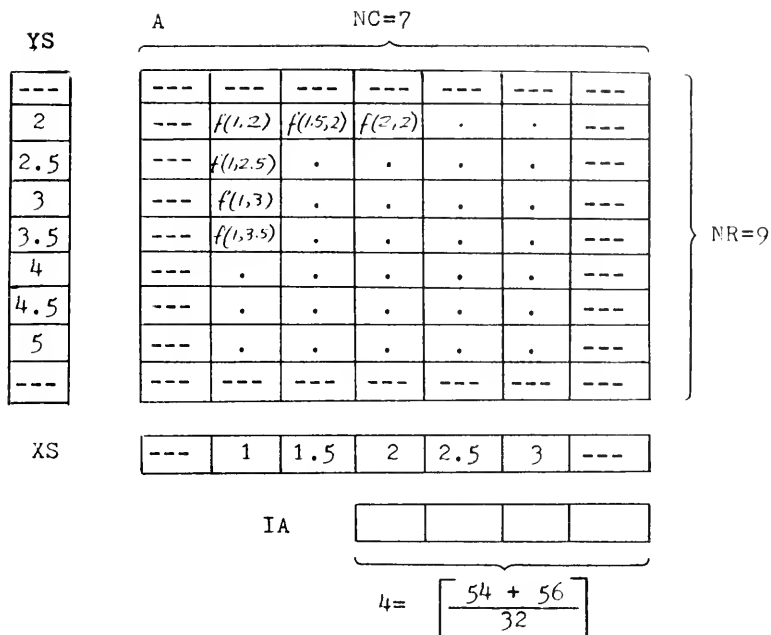
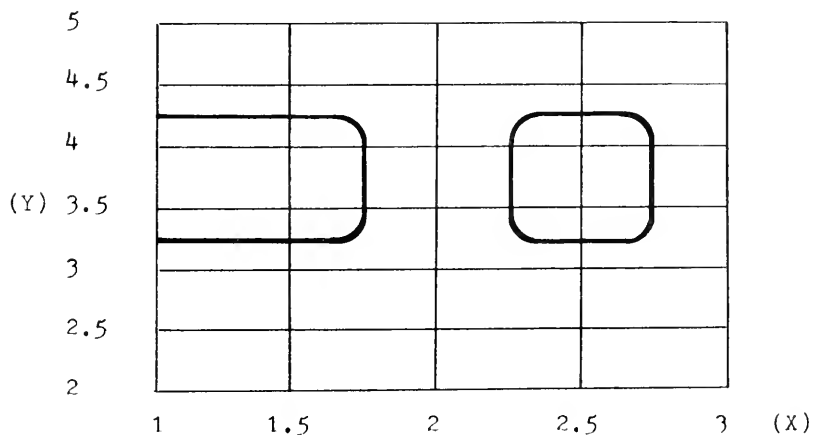


Fig. A.1.- Input variables set up for example.

Fig. A.2.- Grid used in example. $f(x,y)=4.5$ locus' branches.

PTS

NPTS = 15

1.	3.25
1.5	3.25
1.75	2.5
1.75	4.
1.5	4.25
1.	4.25
NA	NA
2.25	3.5
2.5	3.25
2.75	3.5
2.75	4.
2.5	4.25
2.25	4.
2.25	3.5
NA	NA

} NDPTS

Fig. A.3.- Result variables for the locus in Fig. A.2

grid edge. Since in a NC x NR grid there are $(NC-1)*NR + (NR-1)*NC$ edges and one 370 word holds 32 bits, thus the assertion about the dimension of IA in point 2 above.

Bits in IA are manipulated as follows: A bit is set to 1 by means of a call to the subroutine MARK whenever its corresponding grid edge is crossed by the locus to be built. Such circumstance is checked, for every edge but the extra ones, by the test

$$(A1 - ALEVL) * (A2 - ALEVL) \leq 0, \quad (3)$$

where A1 and A2 are the function values at the edge extreme

(3) Refer to section 2 for test details in special cases.

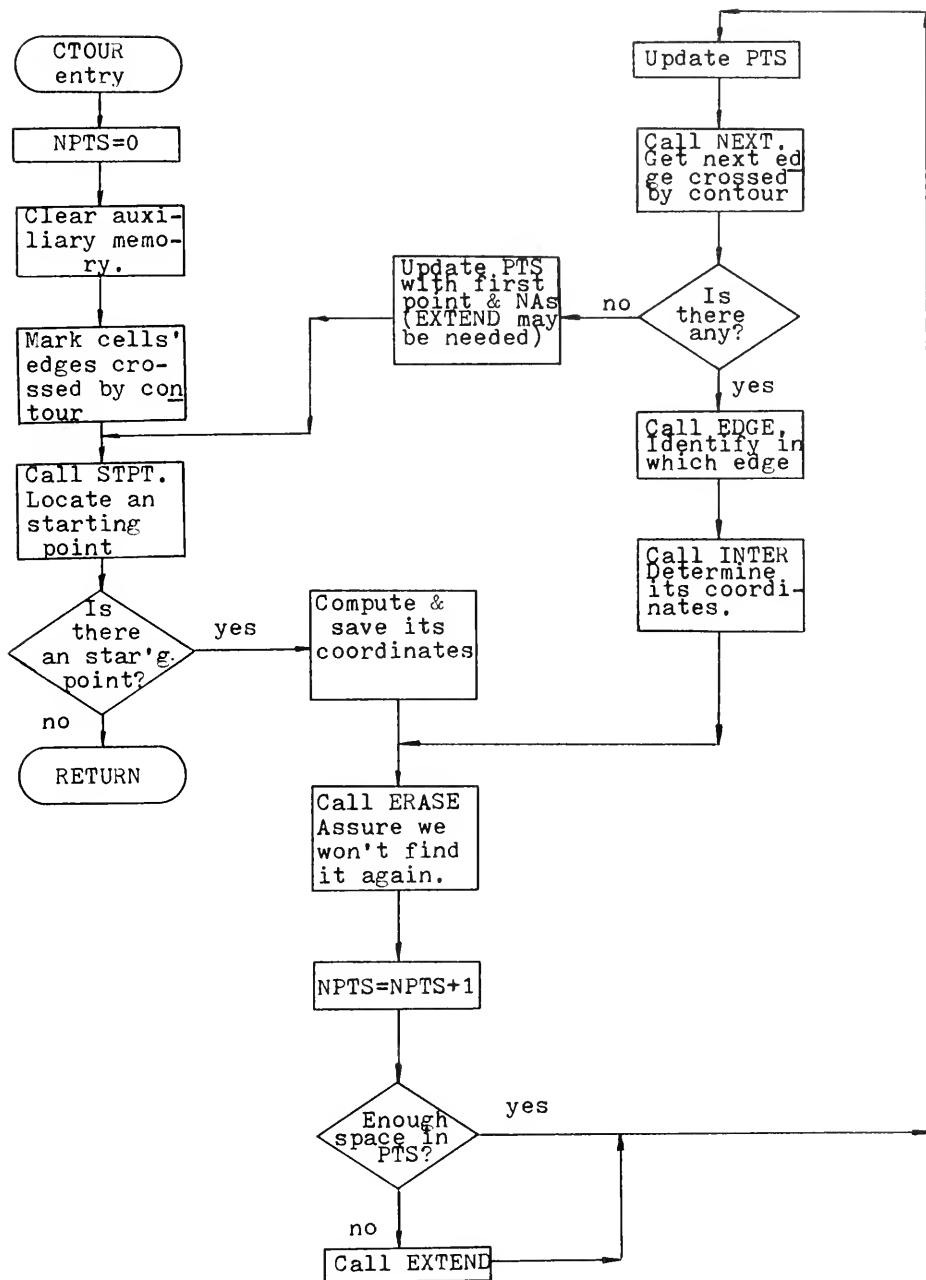


Fig. A.4.- CTOUR flowchart.

points.

The correspondence between IA bits and grid edges is established as follows: Edges are numbered as depicted in Fig. A.5 for the case of a 4 x 7 grid (vertical edges by rows first, then horizontal edges by rows). Having done this, given an edge number there is a corresponding bit in IA. For other uses, however,

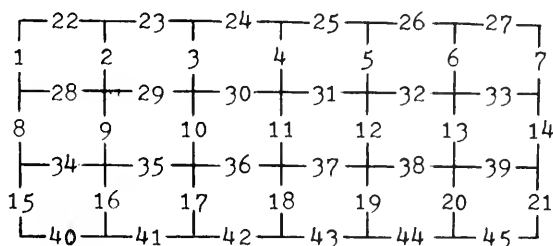


Fig. A.5.- Edge numbers in a 4 x 7 grid.

an edge is more conveniently identified as shown in Fig. A.6: The vertical edge joining the grid points (I,J) and (I+1,J) is speci-

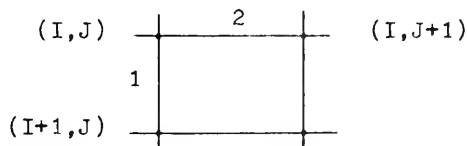


Fig. A.6.- A more convenient edge specification.

fied by means of the 3-tuple (I,J,1), and the horizontal one joining (I,J) to (I,J+1) by (I,J,2). The subroutine INDIA translates such specifications to edge numbers to manipulate the bit memory.

- The subroutine STPT scans the bit indicators in IA looking for 1's by means of calls to the subroutine SEE, which accepts

a bit number and returns a variable = $\begin{Bmatrix} 1 \\ 2 \end{Bmatrix}$ if the bit is $\begin{Bmatrix} a & 0 \\ a & 1 \end{Bmatrix}$.

Boundary edges are checked first, to locate incomplete branches effectively. STPT returns an integer variable K equal to the number of a bit in IA currently set to 1. If no such a bit is found, K is returned set to zero.

- Calling the subroutine EDGE with a bit number K returns an edge specification in the form already seen in Fig. A.6, by means of a 3-tuple (I,J,KIN). (i.e., EDGE performs a translation in edge specification converse to that one done by INDIA).

- INTER interpolates linearly in the edge specified by (I,J,KIN) to end up with a point (XXX,YYY) such that $f(XXX,YYY) = ALEVL$ (with the error proper of a linear interpolation; if such a method is considered ineffective for the function at hand, it is easy to write an alternative INTER subroutine to perform a finer interpolation, although this circumstance is unlikely to arise as other considerations -see section 2- force the grid cells to be of a size where interpolating linearly is usually enough.)

- ERASE accepts a bit number K in IA and sets it to zero.

- NEXT accepts an edge specified as (I,J,KIN) and looks at the adjacent cell searching for another edge crossed by the iso-contour under construction. If such an edge is not found, the original cell itself is checked for a crossed edge (because on the boundary it may happen that the adjacent cell does not have another crossed edge). When this search also fails, the contour branch being treated is complete.

Cells are identified by means of their north - west corner, and their edges by means of a variable set to 1, 2, 3 or 4 according to Fig. A.7. Notice that the 4 edges have to be

identified here in order to effectively locate an adjacent cell.

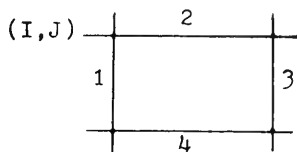


Fig. A.7.- Cell and edge identification for NEXT.

Given a cell as (I,J) and a value $KIN(=1, 2, 3, 4)$ indicating which edge was crossed in that cell, NEXT identifies the adjacent cell as indicated in Fig. A.8 (where x indicates the cross-

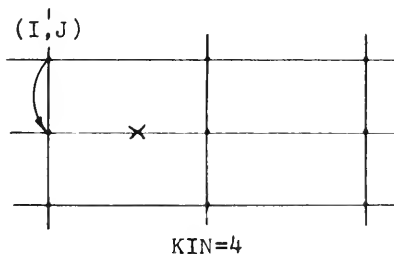
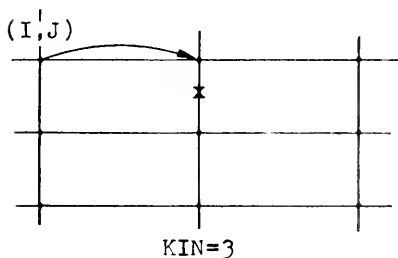
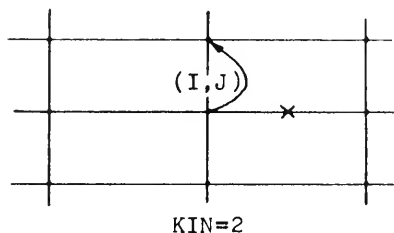
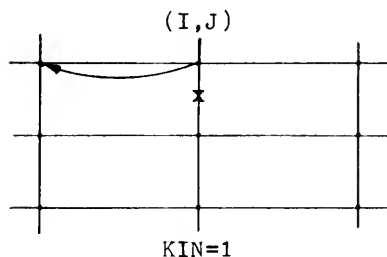


Fig. A.8.- Adjacent cell identification (NEXT).

ssed edge in cell (I,J) and the arrows point to the cell considered adjacent. In addition, NEXT updates I, J and KIN to specify the next edge to be analyzed; it also returns in K the bit number in IA corresponding to it. ($K=0$ indicates no more edges crossed by current locus' branch, thus specifying its end).

10
20
30
40

LEAF 2.0

CHECK

DATE = 75079

09/01/78

```

SUBROUTINE CHECK(IA,I,J,KIN,NR,NC,K)
C****
C ** CHECKS IF THE CELL (I,J) HAS A CROSSING EDGE OTHER THAN KIN,WHICH
C**** WAS ALREADY ERASED. IF SO, RETURNS IT AS (I,J,KIN) AND K. IF NOT,
C**** IT RETURNS K=0.
C****
      DIMENSION IA(1)
      KNEXT=KIN
      DO 30 L=1,3
      KNEXT=KNEXT+1
      IF(KNEXT-4) 20,20,10
10  KNEXT=KNEXT-4
20  K=INDIA(I,J,KNEXT,NR,NC)
      CALL SEE(IA,K,NYES)
      GO TO(30,40),NYES
30  CONTINUE
      K=0
      RETURN
40  KIN=KNEXT
      RETURN
      END

```

CHF00010
 CHF00020
 CHF00030
 CHF00040
 CHF00050
 CHF00060
 CHF00070
 CHF00080
 CHF00090
 CHF00100
 CHF00110
 CHF00120
 CHF00130
 CHF00140
 CHF00150
 CHF00160
 CHF00170
 CHF00180
 CHF00190
 CHF00200
 CHF00210

RELEASE 2.0

C100R

DATE = 75120

15/23/18

C**** HORIZONTAL EDGE CROSSED.

C****

30 K=INDIA(I,J,2,NR,NC)

CALL MARK(IA,K)

40 IF(1-NR*1) 50,70,70

50 IF((A(I,J)-ALFVL)*(A(I+1,J)-ALFVL)) 60,51,70

51 IF(A(I,J)-ALFVL) 60,52,60

52 IF(A(I+1,J)-ALFVL) 60,70,60

C****

C**** VERTICAL EDGE CROSSED.

C****

60 K=INDIA(I,J,1,NR,NC)

CALL MARK(IA,K)

70 CONTINUE

C****

C**** DEFECT A CONTIGUOUS STARTING POINT.

C****

75 CALL STPI(IA,NR,NC,K)

IF(K) 110,110,80

C****

C**** AN STARTING POINT WAS FOUND-IDENTIFY IN WHICH EDGE IS IT.

C****

80 CALL EDGE(K,NR,NC,I,J,KFN)

C****

C**** DETERMINE AND SAVE THE COORDINATES OF THE STARTING POINT TO
C**** BE ABLE TO CLOSE IT WHEN ENDED.

C****

CALL INTER(I,J,KIN,XS,YS,NR,NC,A,ALFVL,XXX,YYY)

FIRSTX=XXX

FIRSTY=YYY

GO TO 86

C****

C**** IDENTIFY CROSSING POINT IN EDGE BY INTERPOLATION.

C****

85 CALL INTER(I,J,KIN,XS,YS,NR,NC,A,ALFVL,XXX,YYY)

C****

C**** FORGET ABOUT THIS CROSSING POINT.

C****

86 CALL ERASE(IA,K)

C****

C**** UPDATE POINTS SET(S).

C****

NPIS=NPIS+1

IF(NPIS=NPIS) 87,90,90

87 CALL EXTEND(PIS,NPIS)

NPIS=1

90 PIS(NPIS,1)=XXX

PIS(NPIS,2)=YYY

C1000490

C1000500

C1000510

C1000520

C1000530

C1000540

C1000550

C1000560

C1000570

C1000580

C1000590

C1000600

C1000610

C1000620

C1000630

C1000640

C1000650

C1000660

C1000670

C1000680

C1000690

C1000700

C1000710

C1000720

C1000730

C1000740

C1000750

C1000760

C1000770

C1000780

C1000790

C1000800

C1000810

C1000820

C1000830

C1000840

C1000850

C1000860

C1000870

C1000880

C1000890

C1000900

C1000910

C1000920

C1000930

C1000940

C1000950

C1000960

1

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RELEASE 2.0

CTHUR

DATE = 75120

15/23/18

```

C***** CTH008700
C*** LOOK FOR NEXT CROSSING POINT. CTH009400
C***** CTH009900
C*** CALL NEXT(IA,IB,IC,I,J,KIN,K) CTH010000
C***** CTH010100
C*** DOES IT EXIST? IF SO, IDENTIFY IT AND CONTINUE CTH010200
C***** CTH010400
C*** IF(K) 100,100,85 CTH010600
C***** CTH010800
C*** IF NOT, TERMINATE A SUBROUT & LOOK IF ANY OTHER EXISTS. CTH011000
C***** CTH011200
100 IF(PTS(NPTS,1)-XS(2)) 101,109,101 CTH011400
101 IF(PTS(NPTS,1)-XS(NCM1)) 102,109,102 CTH011600
102 IF(PTS(NPTS,2)-YS(2)) 103,109,103 CTH011800
103 IF(PTS(NPTS,2)-YS(NRM1)) 104,109,104 CTH012000
104 NPTS=NPTS+1 CTH012200
C*** IF(NPTS=NPTS) 105,106,106 CTH012400
105 CALL EXTEND(PTS,NPTS) CTH012600
NPTS=1 CTH012800
106 PTS(NPTS,1)=FIRSTX CTH013000
PTS(NPTS,2)=FIRSTY CTH013200
109 NPTS=NPTS+1 CTH013400
C*** IF(NPTS=NPTS) 107,108,108 CTH013600
107 CALL EXTEND(PTS,NPTS) CTH013800
NPTS=1 CTH014000
108 PTS(NPTS,1)=ZNA CTH014200
PTS(NPTS,2)=ZNA CTH014400
GO TO 75. CTH014600
C***** CTH014800
C*** TERMINATE. RETURN TO CALLING PROGRAM. CTH015000
C***** CTH015200
110 RETURN CTH015400
END CTH015600

```


RELEASE 2.0

EDGE

DATE = 75084

10/31/07

SUBROUTINE EDGE(K,NR,NC,I,J,KIN)

FDG00010

FDG00020

FDG00030

FDG00040

FDG00050

FDG00060

FDG00070

FDG00080

FDG00090

FDG00100

FDG00110

FDG00120

FDG00130

FDG00140

FDG00150

C****

C ** TRANSLATES K-> I,J,KIN

C****

KK=K-((NR-1)*NC)

IF(KK) 10,10,20

10 KIN=1

I=((K-1)/NC)+1

J=K-((I-1)*NC)

RETURN

20 I=((KK-1)/(NC-1))+1

J=KK-((I-1)*(NC-1))

KIN=2

RETURN

END

LEASE 2.0

INDIA

DATE = 75083

11/34/16

FUNCTION INDIA(I,J,KIN,NR,NC)

IND00010

IND00020

IND00030

IND00040

IND00050

IND00060

IND00070

IND00080

IND00090

IND00100

IND00110

IND00120

IND00130

IND00140

IND00150

IND00160

IND00170

IND00180

IND00190

IND00200

C****

C ** COMPUTES THE INDEX IN IA CORRESPONDING TO THE EDGE SPECIFIED

C **** BY I, J & KIN. NR & NC ARE NO. OF ROWS & COLUMNS OF A AS DECLARED

C **** IN THE CALLING PROGRAM (INCLUDING MANDATORY EXTRA ROWS AND COLS.)

C****

II=I

JJ=J

KKIN=KIN

GO TO(40,30,10,20),KKIN

10 JJ=JJ+1

KKIN=1

GO TO 40

20 II=II+1

KKIN=2

30 INDIA=(NR-1)*NC+(NC-1)*(II-1)+JJ

RETURN

40 INDIA=JJ+(II-1)*NC

RETURN

END

LEASE 2.0

INTER

DATE = 75079

16/00/42

```

SUBROUTINE INTER(I,J,KIN,XS,YS,NR,NC,A,ALFVL,XXX,YYY)
C****
C  ** A CONTIUR POINT IS GENERATED VIA LINEAR INTERPOLATION, LYING IN THE
C**** GRID EDGE DEFINED BY (I,J,KIN). THE RESULT IS (XXX,YYY).
C****
      DIMENSION XS(NC),YS(NR),A(NR,NC)
      II=I
      JJ=J
      GO TO(40,30,10,20),KIN
10  JJ=JJ+1
      GO TO 40
20  II=II+1
30  YYY=YS(II)
      DELT=XS(JJ+1)-XS(JJ)
      P=(A(II,JJ+1)-A(II,JJ))/DELT
      XXX=(ALFVL-A(II,JJ))/P
      XXX=XXX+XS(JJ)
      RETURN
40  XXX=XS(JJ)
      DELT=YS(II+1)-YS(II)
      P=(A(II+1,JJ)-A(II,JJ))/DELT
      YYY=(ALFVL-A(II,JJ))/P
      YYY=YYY+YS(II)
      RETURN
END

```

100

101

102

103

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107

108

109

110

SOURCE STATEMENT

ASM 0105 21.14 03/25/77

1	MARK	START		MAR0001
2		ENTRY FRASE,SEF		MAR0002
3		STM 14,12,12(13)		MAR0003
4		HALR BASE,0		MAR0004
5		USING *,BASE		MAR0005
6	AC	EQU 8		MAR0006
7	ADDR	EQU 9		MAR0007
8	BASE	EQU 10		MAR0008
9	INDEX	EQU 11		MAR0009
10	RUNE	EQU 2		MAR0010
11	MASK	EQU 3		MAR0011
12	K	EQU 4		MAR0012
13		L RUNE,=F'1'		MAR0013
14		LR MASK,RUNE		MAR0014
15		SLL MASK,31		MAR0015
16		L ADDR,4(0,1)	*GET K.	MAR0016
17		L K,0(0,ADDR)	*	MAR0017
18		LR AC,K	.	MAR0018
19		SR AC,RUNE	.	MAR0019
20		SRL AC,5	.GET CORRESPONDING ARRAY ELEMENT	MAR0020
21		SLL AC,2	.	MAR0021
22		LR INDEX,AC	.	MAR0022
23		SLL AC,3	*	MAR0023
24		LR 6,AC	*	MAR0024
25		LR AC,K	*GET HIT POSITION	MAR0025
26		SR AC,6	*	MAR0026
27		SR AC,RUNE	.	MAR0027
28		SRL MASK,0(AC)	.PREPARE MASK	MAR0028
29		L ADDR,0(0,1)	*	MAR0029
30		L AC,0(INDEX,ADDR)	*GET AFFECTED ARRAY WORD	MAR0030
31		LR AC,MASK	MARK SPECIFIED BIT	MAR0031
32		ST AC,0(INDEX,ADDR)	STORE WORD BACK	MAR0032
33		LM 14,12,12(13)	*RETURN	MAR0033
34		RR 14	*	MAR0034
35	FRASE	STM 14,12,12(13)		MAR0035
36		HALR BASE,0		MAR0036
37		USING *,BASE		MAR0037
38		L RUNE,=F'1'		MAR0038
39		LR MASK,RUNE		MAR0039
40		SLL MASK,31		MAR0040
41		L ADDR,4(0,1)	*GET K	MAR0041
42		L K,0(0,ADDR)	*	MAR0042
43		LR AC,K	.	MAR0043
44		SR AC,RUNE	.	MAR0044
45		SRL AC,5	.GET CORRESPONDING ARRAY ELEMENT	MAR0045
46		SLL AC,2	.	MAR0046
47		LR INDEX,AC	.	MAR0047
48		SLL AC,3	*	MAR0048
49		LR 6,AC	*	MAR0049
50		LR AC,K	*GET HIT POSITION	MAR0050
51		SR AC,6	*	MAR0051
52		SR AC,RUNE	.PREPARE MASK	MAR0052
53		SRL MASK,0(AC)	.	MAR0053
54		L ADDR,0(0,1)	*GET AFFECTED ARRAY WORD	MAR0054
55		L AC,0(INDEX,ADDR)	*	MAR0055

SOURCE STATEMENT

ASM 0105 21.14 03/25/77

66	XR	AC, MASK	ERASE SPECIFIED BIT	MAR0056
67	ST	AC, 0 (INDEX, ADDR)	STORE WORD BACK	MAR0057
68	LM	14, 12, 12 (13)	*RETURN	MAR0058
69	RR	14	*	MAR0059
70	ST	STM 14, 12, 12 (13)		MAR0060
71	BALR	BASE, 0		MAR0061
72	USING	*, BASE		MAR0062
73	L	RONE, =F'1'		MAR0063
74	LR	MASK, RONE		MAR0064
75	SLL	MASK, 31		MAR0065
76	L	ADDR, 4 (0, 1)	*GET K	MAR0066
77	L	K, 0 (0, ADDR)	*	MAR0067
78	LR	AC, K	.	MAR0068
79	SR	AC, RONE	.	MAR0069
80	SRL	AC, 5	.GET CORRESPONDING ARRAY ELEMENT	MAR0070
81	SLL	AC, 2	.	MAR0071
82	LR	INDEX, AC	.	MAR0072
83	SLL	AC, 4	*	MAR0073
84	LR	6, AC	*	MAR0074
85	LR	AC, K	*GET BIT POSITION	MAR0075
86	SR	AC, 6	*	MAR0076
87	SR	AC, RONE	.PREPARE MASK	MAR0077
88	SRL	MASK, 0 (AC)	.	MAR0078
89	L	ADDR, 0 (0, 1)	*GET AFFECTED ARRAY WORD.	MAR0079
90	L	5, 0 (INDEX, ADDR)	*	MAR0080
91	L	ADDR, 8 (0, 1)	GET NYYES ADDRESS	MAR0081
92	NR	MASK, 5	.CHECK WHETHER BIT WAS ON OR NOT	MAR0082
93	BC	4, YES	.	MAR0083
94	ST	RONE, 0 (0, ADDR)	*IF NOT, SET NYYES TO 1	MAR0084
95	B	OUT	*AND RETURN	MAR0085
96	SLL	RONE, 1	.IF YES, SET NYYES TO 2	MAR0086
97	ST	RONE, 0 (0, ADDR)	.	MAR0087
98	LM	14, 12, 12 (13)	*RETURN	MAR0088
99	RR	14	*	MAR0089
100	END			MAR0090
101		=F'1'		

EASE 2.0

NEXI

DATE = 75079

15/44/11

SUBROUTINE NEXI(IA,NR,NC,I,J,KIN,K)

NEX00010

C****

NEX00020

C * LOOKS FOR NEXT CROSSING EDGE, IF ANY & RETURNS IT IN BOTH FORMS

NEX00030

C**** (I,J,KIN) & K.-REFERRING TO IA-/RETURNS K=0 IF IT DOESN'T EXIST

NEX00040

C****

NEX00050

DIMENSION IA(1)

NEX00060

II=I

NEX00070

JJ=J

NEX00080

C****

NEX00090

C****

EDGE OF ADJACENT CELL

NEX00100

C****

NEX00110

KKIN=KIN+2

NEX00120

IF(KKIN-4) 20,20,10

NEX00130

10 KKIN=KKIN-4

NEX00140

20 GO TO(30,40,50,60),KIN

NEX00150

30 JJ=JJ-1

NEX00160

GO TO 70

NEX00170

40 II=II-1

NEX00180

GO TO 70

NEX00190

50 JJ=JJ+1

NEX00200

GO TO 70

NEX00210

60 II=II+1

NEX00220

70 CALL CHECK(IA,II,JJ,KKIN,NR,NC,K)

NEX00230

IF(K) 80,80,90

NEX00240

C****

NEX00250

C****

WHEN IT DOESN'T EXIST IN THE ADJACENT CELL,CHECK THE CURRENT ONE.

NEX00260

C****

NEX00270

C 30 CALL CHECK(IA,I,J,KIN,NR,NC,K)

NEX00280

RETURN

NEX00290

90 I=II

NEX00300

J=JJ

NEX00310

KIN=KKIN

NEX00320

RETURN

NEX00330

END)

NEX00340

- 2,

√10

SC, SC

, K²

(1) (1)

(10), (10)

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LEASE 2.0

SEARCH

DATE = 75086

09/23/04

SUBROUTINE SEARCH(A,NR,NC,ALFVL,NCNTR)

SF000010

SF000020

** SCANS THE TABLE A AND GENERATES NCNTR VALUES EQUALLY SPACED

SF000030

**** BETWEEN THE MAXIMUM AND MINIMUM VALUES OF A.

SF000040

SF000050

DIMENSION A(NR,NC),ALFVL(NCNTR)

SF000060

AMAX=A(1,1)

SF000070

AMIN=AMAX

SF000080

NRM1=NR-1

SF000090

NCM1=NC-1

SF000100

DO 40 I=2,NRM1

SF000110

DO 40 J=2,NCM1

SF000120

IF(A(I,J)-AMAX) 20,20,10

SF000130

10 AMAX=A(I,J)

SF000140

GO TO 40

SF000150

20 IF(A(I,J)-AMIN) 30,40,40

SF000160

30 AMIN=A(I,J)

SF000170

40 CONTINUE

SF000180

J=NCNTR-1

SF000190

DELTA=(AMAX-AMIN)/J

SF000200

ALFVL(1)=AMIN

SF000210

DO 50 I=1,J

SF000220

50 ALFVL(I+1)=ALFVL(I)+DELTA

SF000230

RETURN

SF000240

END

SF000250

2.0

SUBR

1

1.1 DIMENSION

1.2 DIMENSION

1.3 DIMENSION

1.4 DIMENSION

1.5 DIMENSION

1.6 DIMENSION

1.7 DIMENSION

1.8 DIMENSION

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1.54 DIMENSION

1.55 DIMENSION

LEAST 2.0

STPT

DATE = 75084

14/00/17

SUBROUTINE STPT(IA,NR,NC,K)

```

****
** LOOKS FOR AN EDGE MARKED IN IA.IT GIVES PREFERENCE TO BOUNDARY
**** EDGES,SO THAT INCOMPLETE SUBTOURS ARE LOCATED FIRST AND THEIR
**** STARTING POINTS ARE ON THE BOUNDARIES.
****

```

DIMENSION IA(1)

NRM1=NR-1

NCM1=NC-1

I=2

5 DO 10 J=2,NCM1

K=INDIA(I,J,2,NR,NC)

CALL SEE(IA,K,NOYES)

GO TO(10,90),NOYES

10 CONTINUE

IF(I-2) 20,20,30

20 I=NRM1

GO TO 5

30 J=2

40 DO 50 I=2,NRM1

K=INDIA(I,J,1,NR,NC)

CALL SEE(IA,K,NOYES)

GO TO(50,90),NOYES

50 CONTINUE

IF(J-2) 60,60,70

60 J=NCM1

GO TO 40

70 NRM2=NRM1-1

NCM2=NCM1-1

DO 80 I=2,NRM2

DO 80 J=2,NCM2

DO 80 L=1,2

K=INDIA(I,J,L,NR,NC)

CALL SEE(IA,K,NOYES)

GO TO(80,90),NOYES

80 CONTINUE

K=0

90 RETURN

END

```

STP00010
STP00020
STP00030
STP00040
STP00050
STP00060
STP00070
STP00080
STP00090
STP00100
STP00110
STP00120
STP00130
STP00140
STP00150
STP00160
STP00170
STP00180
STP00190
STP00200
STP00210
STP00220
STP00230
STP00240
STP00250
STP00260
STP00270
STP00280
STP00290
STP00300
STP00310
STP00320
STP00330
STP00340
STP00350
STP00360
STP00370
STP00380
STP00390

```


APPENDIX BRoutine invocation from TROLL. Macros.

The two TROLL macros listed below facilitate invocations to CTOUR from the TROLL system environment. (See TROLL manuals for details in specific commands*). To invoke these macros from any user account, it is necessary to specify the following command after logging in:

```
SEARCH MIT128_MACRO;
```

The macro &LIKPLOTS includes computation of likelihood function values at grid points, so that it is only useful for our particular problem. &PLOTS is more general, but it requires the function table to be available prior to its invocation.

Calling sequences examples are also included and are self-explanatory. (Capital letters correspond to system prompts).

* Available from the National Bureau of Economic Research, 545 Technology Sq., Cambridge.



UNREPLACES

```

ERROR: EGOTO ERR-
SEARCH WILL FUNCTION:
OUTPUT DEVICE T4010:
OUTPUT DIMENSIONS:
OUTPUT EOG(IGR) OF OR 'MIGR10':":
READ E4"YOUR MIN VALUES:" &END
READ E5"YOUR SIGMA SQ. VALUES:" &END
DO R44=EXPAND(E4,NOB(E5),NOB(E5)):DO R45=TRANSP(R44):
DO R55=EXPAND(E5,NOB(E4),NOB(E4)):
READ E7"R VALUES:" &END
READ E8"OBSERVATIONS:" &END
READ E9"R VALUE:" &END
DO E8E9=1 IKTOT(E7,R44,R45,COUNTIME(E9),E8):
DO T=ARG(1)=MAXS(E8E9):
PRINTA
YOUR MAXIMUM VALUE IS E1FEARG(1).
CHOOSE YOUR CONTINUR LEVELS ACCORDINGLY.
&END
READ E2"NEW LEVELS?" &END
E1F E2 C=0 NO EGOTO DLD E1FEEND
DEFDEF DATA E3"NEW LEVELS NAME:"":
SIGMA:
DEFIT E3,1,1:
ADD TOP,CONTINUR LEVELS IN QUITES:"":
FILE:
EGOTO NEXT
EOLD:
READ E3"OLD LEVELS NAME:" &END
ENEXT:
DO E3=REV(SORT(E3,E3)):
DO SETFEARG(E3,50):
DO E8E9C=NEBCIRRE(E4,E5,E8E9,E3):
DO E4.=E8E9C_X:
DO E5.=E8E9C_Y:
DO T=ARG(2)=NOB(E8E9C):
DO SETFEARG(E8E9C,5):
ESET E1FEARG(3)=1 &END
ESET E1FEARG(4)=5 &END
ESETC E1FEARG(1)=ABCDEFGHIJKLMNPQRSTUWXYZ &END
CLIMDS:
ESET E1FEARG(1)=1 &END
DESPACE E4. E5.:
READ E"OVARIALE IN HORIZONTAL AXIS:" &END
E1F E0 C=0 "E5" EGOTO REMVS &FEEND
PLANE 1 2:
EGOTO DIBOX
REMOVES:
PLANE 2 1:
DIBOX:
ESETC E1FEARG(2)=EKEEPL 1 E1FEARG(1) &END
ESETC E1FEARG(1)=E5IRIP 1 E1FEARG(1) &END
MARK E1FEARG(1) E1FEARG(2):
ESET E1FEARG(1)=E1FEARG(E1FEARG(4))+1 &END
E1F E1FEARG(3) 1,1 E1FEARG(2)
ESET E1FEARG(3)=E1FEARG(3)+1 &END
ESET E1FEARG(4)=E1FEARG(4)+1 &END
EGOTO DIBOX
&FEEND
SCALE:
PRINTA
Z/GO/ZE/RSI 3 CONTINUR VALUES ARE (THE PRES. OF T4010 4):

```


$A \equiv \exists x [x \in A \wedge (x \in \{ \emptyset \})]$, $B \equiv \exists x [x \in A \wedge (x \in \{ \emptyset \})]$, $C \equiv \exists x [x \in A \wedge (x \in \{ \emptyset \})]$.

$\exists x (x \in \{ \emptyset \})$

$\exists x (x \in \{ \emptyset \})$

$\exists x (x \in \{ \emptyset \})$

$\exists x [x \in \{ \emptyset \} \wedge (x \in \{ \emptyset \})] \equiv \exists x [x \in \{ \emptyset \} \wedge (x \in \{ \emptyset \})]$

$\exists x (x \in \{ \emptyset \}) \equiv \exists x (x \in \{ \emptyset \})$

TROLL COMMAND: do mu = seq(0.,5.,0.1);

TROLL COMMAND: do sigma = seq(0.01,5.,0.1);

TROLL COMMAND: &likplots

'GRID' OR 'NOGRID':grid

YOUR MU VALUES:mu

YOUR SIGMA SQ. VALUES:sigma

B VALUES:bb

OBSERVATIONS:nsob

N VALUE:300

%YOUR MAXIMUM VALUE IS -128.125.
CHOOSE YOUR CONTOUR LEVELS ACCORDINGLY.
NEW LEVELS?yes

NEW LEVELS NAME:lev300

%NEW SERIES LEV300

ENTER LEVELS IN QUOTES:"-128.125 -128.5 -128.0 -129 -130 -135"

VARIABLE IN HORIZONTAL AXIS:sigma

TROLL COMMAND: &plots

'GRID' OR 'NOGRID':grid

YOUR TABLE:nsob300

THE MAXIMUM VALUE IN NSOB300 IS -128.125.
CHOOSE YOUR CONTOUR LEVELS ACCORDINGLY.
NEW LEVELS?no

OLD LEVELS NAME:lev300

YOUR X VALUES:mu

YOUR Y VALUES:sigma

YOUR CONTOUR NAME:nsob300c

%VARIABLE IN HORIZONTAL AXIS:sigma

APPENDIX CProgram requirements and performance.

The memory requirements for the subroutines included in the iso - contour generating package are as follows:

<u>Subroutine</u>	<u>Bytes</u>
CHECK	684
CTOUR	2,552
EDGE	604
INDIA	628
INTER	1,134
MARK, ERASE, SEE	240
NEXT	822
SEARCH	896
STPT	1,088
Total	<hr/> 8,648

As for execution times, we include below the CPU times taken to generate and plot iso - contours in different cases. We measured them for different grid sizes (given by the number of grid nodes) and different number of contours to generate. Actually, these times are execution times for the macro &PLOTS, and thus they include some TROLL overhead. They are measured in seconds.

<u>Grid Size</u>	<u>Contours generated and plotted</u>	
	<u>3</u>	<u>6</u>
25 x 25	0.67	0.92
50 x 50	1.42	2.34
75 x 75	2.43	4.28
100 x 100	3.82	6.91

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